### IMAGE COMPRESSION USING WAVELET TRASFORM

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### **Certificate**

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 $\operatorname{in}$ 

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# Dedicated to my Parents, Brothers and Sisters

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*Muhammad Shoaib* 

#### **ABSTRACT**

A wavelet based algorithm for digital image compression is presented. The algorithm has good compression ratio as well as PSNR. A comparison with the standard JPEG algorithm for a number of images is also presented. This algorithm has been applied to gray scale images but can be extended for color.

## **Contents**





## **Chapter 1**

## **Introduction**

This work deals with the image coding using wavelet transform. It can be better described by casting it in the framework of transform image coding . In this typical state. of the art, transform image coding system, the encoder consists of a linear transform operation, followed by quantization of the transform-domain coefficients, and lossless compression of the quantized coefficients using an entropy coder. After the encoded bit. stream of an input image is transmitted over the channel (assumed to be perfect.), the decoder undoes all the functionalities applied in the encoder and tries to reconstruct a decoded image that looks as close as possible to the original input image, based on the transmitted information. This source coding paradigm has become the de-facto standard for lossy image compression application such as JPEG  $[1]$  and MPEG  $[2]$ , where the only loss of information occurs in the quantizer. See Figure  $1.1$ .

The basic idea behind using a linear transformation is to make the task of compressing an image in the transform domain after quantization easier than direct coding in the spatial domain. A good candidate transformation should be able to offer flexible image representation decorrelation ( to facilitate efficient entropy coding) and good energy compaction in the transform domain (so that fewer quantized coefficients are needed to be encoded and the rest can be discarded for minimum distortion). It is also desirable for the transform to be orthogonal so that the energy is conserved from the spatial

domain to the transform domain , and the distortion in the spatial domain introduced by quantization of transform coefficients can be directly examined in the transform domain . Finding the optimal orthogonal transform for an  $N \times N$  image necessitates a search over the set of all  $N^2 \times N^2$  unitary matrices, which is clearly impossible, since such a set. of unitary matrices is infinite. In practice, suboptimal approximations such as the discrete cosine transform (DCT) are used for computational efficiency and being image independent [3] .



Figure 1.1: Block diagram of typical transform coding system (a ) The encoder block (b) The decoder block diagram

Transform coding (DCT) is an efficient block oriented image compression techniques that is now being widely used in the image compression industry. while various discrete transforms have been investigated for application to transform coding, only the discrete.  $\cos$ ine transform (DCT) has emerged as the most practical and efficient transform [11]. The principle used for transform coding is to resolve the original subpicture into a linear combination of a set of predefined subpictures, called *basis function*. The transform coefficients, which form the transmitted informations, are the multiplying factors of the basis functions. The bit rate reduction is operated by quantizing the transform coefficients before transmission.

To achieve a high compression ratio, most of the transform coefficients are coarsely quantized. Coarse quantization of the transform coefficients results in various artifacts in the coded images. Methods for reducing the block-artifacts in the transform coding of the images has been extensively studied. One of them is to incorporate the HVS (Human Visual System) properties with designing the quantization matrix (Q-matrix) for the transform coding since the human eye has discriminative sensitivity to different. spatial frequencies. Another approach to cope with this problem is to replace the DCT basis functions by discrete wavelet basis [4], which has shorter basis functions for higher frequencies, and longer basis functions for lower frequencies as shown in figure [3.1]. There are more samples to represent the higher frequency sub-bands, than the lower frequencies ones. Therefore, sharp edges, which are well localized spatially and have significant highfrequency contents, can be represented more compactly with the DWT than with the DCT.

Recently, wavelets and filter bank theory  $[5, 6, 7, 8, 9, 10, ]$ , together with their generalization such as wavelet packet ( $W P$ ) [5, 6, 20,], have appeared as alternatives to the DCT basis due to their ability to provide more flexible space-frequency resolution trade-off image representation.

#### **1.1 Layout of the dissertation**

The research work carried out and described in this dissertation constitutes a lossy compression algorithm. it is based on wavelet transform. Before going into details of the actual work existing image compression methods and wavelet based algorithm are discussed. chapter 2 discusses the need for image compression and standard image compression techniques. Chapter 3 deals with the wavelet analysis. Compression using wavelet transform and its results are discussed in chapter 4. Some real images along with the processed images are presented in chapter 4.

## **Chapter 2**

## **Image Compression**

### **2 .1 Introduction**

Digital image contain large amounts of information therefore from the stand points of data storage and transmission, one would like to have an image stored in a way that. requires fewer bits. If one is processing a large number of images, then efficient representation is necessary in order not to overwhelm memory and if images are to be transmitted then there are bandwidth limitations so that timely transmission requires efficient representation. Thus data need to be compressed, or coded in such a way as to facilitate storage and transmission. Further more, various image representations enhance the speed of various algorithms. A key here is elimination of redundancy, and various transformation serve to reduce various type of redundancy. More than simply finding an efficient image representation, image are often altered in a noninvertible manner. Since this involve the loss of information, such compression is termed 'lossy' as opposed to invertible encoding, which are termed lossless.

### **2.2 Why compression is necessary?**

"A picture is worth a thousand words." This English aphorism reminds us of the importance of images. It is especially true in the age of information highway and multimedia. Computers, fax machines, video phones, teleconferencing system and storage devices impact our workplace. Text, data, sound, images and video clip are grouped together to send over data networks or to store The amount of data is astronomical. Compression increases the throughput of the network and the capacity of the storage device .For satellite transmission, compression greatly reduces cost.

To understand why compression is needed it is a good idea to start with the analog video signal. The analog video signal is mostly the source for a digital video, because the most common form of the video signal in nse today

The full resolution of PAL video signal is 720\*576=414,720 pixels for one picture . For true color 24 Bites per pixel are needed, so  $720*576*24/8=1,244,160$ . bytes are needed for one picture. The PAL video signal contains 25 pictures per second. so we need 31,104,000 bytes per second digital video. This means we have bit rate of 248,832,000 Bytes per second or about 249Mbit/sec.

If you now compute how much space you need to store one 90 minute movie you will recognize why compression is needed:

90 Minutes contain  $90*60=5400$  seconds. One 90 minute movie needs  $5400*31,104,000=$ 16,796,160,000 bytes or 156 G Byte.

This astronomical bit rate can not be handled by any computer system today. The logical solution to this problem is digital image compression.

### **2.3** Lossless And Lossy Compression

Compression techniques are classified into two categories lossless and lossy as shown in figure  $[2.1]$ . Lossless techniques are capable to recover the original data perfectly. These algorithms are used to archive computer data which have to recovered perfectly.



Figure 2.1: Basic element of lossey and lossless image compression

Lossy techniques involve algorithms which recover only similar data to the original one. The lossy techniques provide higher compression ratios, and therefore they are more often applied in the image compression than lossless techniques.

The lossy compression algorithm make use of the characteristic of the human eye, because a picture contains some information which is not necessary for the pictnre quality. The human visual system does not treat all the visnal information with equal sensitivity For example, the eye is more sensitive for changes in the luminance than in the chrominance. For this reason it suffies to transfer only one chrominance pixel U and V for four luminance pixel Y. This indicates that four adjacent pixels have the same color information but various brightness on the display screen. These make it possible to reduce the color information .

The human eye is also less sensitive to high frequencies. So it is a good idea to transfer the lower frequencies more exactly than the high frequencies or to clip the high frequencies. This reduction is not exactly reversible . But because of the human eye there are some possibilities to cut high frequencies and to reduce the color information without visible artifacts.

#### **2.4 JPEG**

The first standard for image compression was developed by the JPEG (Joint Photographic Experts Group) committee in the eighties.

The JPEG committee was introduced by the ISO / IEC ( International Standard Organization / International Electro technical Commission). The JPEG compression algorithm is used for still image applications. The JPEG standard is targeted for fullcolor still frames, achieving 15:1 average compression. You can get a compression factor up to 30 nearly without visible quality difference. The JPEG standard provides four modes of operation:

- 1. *sequential DCT based encoding* in which each image component is encoded in a single left-to-right and top-to bottom scan.
- *2. progressive DCT based encoding* in the image is encoded in multiple scan, in order to produce a quick rough decoded image when the transmission time is long. This encoding mode is often used in the World Wide Web to provide a quick transmitted picture which gets better and better.
- *3. lossless encoding* in which the image is encoded to guarantee the exact reproduction. But the compression ration is only about 2:1
- *4. hierarchical encoding* in which the image is encoded in multiple resolution.

The JPEG encoding is based on the transformation in the frequency domain using

the Discrete Cosine Transform  $($  DCT  $)$ . The compression is made by a quantization in the frequency domain and this means that JPEG is a lossy compression technique.

### 2.4.1 Transformation From The Spatial Domain To The Fre**quency D omain**

The standardized image compression techniques transform a spatial domain into the frequency domain. This transformation into the frequency domain exploits the spatial correlation of the pixels by converting them to a set of independent coefficients. This transformation is done on a block basis. The picture is divided mostly into 8\*8 pixel blocks and then these blocks are transformed using a 2-D transformation into the frequency domain.

The video compression algorithms use the discrete cosine transform (DCT). The DCT offers the advantage in comparison with other transformations, that the coefficients are all in real domain. There are other transformations available, but in the JPEG and MPEG standard the transformation is done by the DCT.

The forward 2-D DCT and the inverse 2-D IDCT are defined as follows: DCT

$$
F(u, v) = \frac{1}{4} * C(u) * C(v) * \sum_{j=k}^{n} \sum_{k}^{n} f(j, k) * \left[ \cos \frac{(2+j+1) * u + \pi}{16} \right] * \left[ \cos \frac{(2*k+1) * u + \pi}{16} \right]
$$
  
IDCT  

$$
F(J, k) = \frac{1}{4} \sum_{j=k}^{n} \sum_{k}^{n} C(u) * C(v) * \left[ \cos \frac{(2+j+1) * u + \pi}{16} \right] * \left[ \cos \frac{(2*k+1) * u + \pi}{16} \right]
$$
  
with  

$$
C(w) = \left\{ \frac{1}{\sqrt{2}} \quad with \quad w = 0 \right\}
$$

The result of the DCT is an  $8 \times 8$  matrix with the frequency coefficient. In the value at the upper left corner is the DC coefficient. The frequency increases to the right and to the bottom.

Normally a picture contain many low frequency components. Mostly the coefficients are concentrated in the upper left corner of a matrix.

After the transformation into the frequency domain the results of the DCT are quantized. The higher frequencies are more quantized than the lower frequencies, because the human eye is not very sensitive for the higher frequency components. The result of the quantization of higher frequencies is mostly zero. Because the high frequency coefficient are more quantized as the low ones.

#### **2.5 Quantization**

A quantizer is essentially a staircase function that maps the possible input values into a smaller number of output levels. In this way the number of symbols that need to be encoded are reduced at the expense of introducing error in the reconstructed image. The type and degree of quantization has a large impact on the final bit rate and the reconstructed picture quality of a lossy scheme. The individual quantization of each signal value is called scalar quantization ( SQ ), and the joint quantization of a block of signal value is called vector quantization ( VQ ). The selection of a quantizer is usually based on the minimization of some distortion measure for a given average output bit rate.

#### **2.5.1 Scalar Quantization**

As mentioned previously, scalar quantization ( SQ ) refers to the independent quantization of each signal value. The main advantage of SQ is its implementation simplicity and optimal performance in many situation.

#### **2.5.2 Vector Quantization**

The joint quantization of a block of signal value is called vector quantization (VQ ). In VQ, an N-dimentional input vector  $X=[x_1,x_2, \ldots, x_n]$ , whose components represent the discrete or continuous signal value, is mapped into one of *N* possible reconstruction

vectors  $Y_i$ , donated by  $d [X, Y_i]$  and is defined according to the application. The most common distortion measurement is mean-squared (MSE), given by

$$
d_{mse}(X, Y) = \frac{1}{n} \sum (x_i - y_i)^2
$$

The set Y is sometimes referred to as the reconstruction *code book* and its members are called *codevector s* or *templates.* The *codebook design* problem is to find the optimal code book ( in the sense of minimizing average distortion ) for given input signal statistics, distortion measure, and code book size, N.

#### **Hierarchical VQ**

Compute the variance of an  $(N \times N)$  block. If it is above a given threshold, divide it into four( $N/2 \times N/2$ ) block. Compute the variances of the four  $(N/2 \times N/2)$  blocks. If all or some of the variances of the  $(N/2 \times N/2)$  blocks are above an other threshold, then divide each  $(N/2 \times N/2)$  block (that is, those above the threshold) into four  $(N/4 \times N/4)$ blocks. This is a variable block size and, hence, a variable vector dimension VQ reflecting the image activity. Separate code book for each block size are to be designed and, of course, stored at the encoder and decoder. Overhead by bits indicating the actual coding mode need to be transmitted. This adaptivity improves the coding efficiency at the cost of increased complexity.

### **2.6 Symbol Encoding**

The last component in common compression algorithms is symbol encoding. i.e., mapping of output symbols (values) resulting from the decomposition and / or quantization stages into channel symbols. The mapping operation is also referred to as noiseless coding, lossless coding or data compaction coding. Symbol encoding may be as simple as using fixed-length binary code words to represent symbols, or it might us variable-length code • words for better efficiency.

#### **2.6.1 Fixed Length Codes**

In some application it is desirable to use fixed length code words to minimize implementation complexity or to satisfy certain channel constraints, such as the need for a constant data rate. In fixed length coding, each source symbol is assigned a fixed length code, where code length depends upon the number of symbols the source can generate. The inefficiency of this type of coding results when the number of symbols is not a power of 2. The variable-length coding is better choice in such cases.

#### **2.6.2 Variable - Length Codes**

Variable-length codes are efficient than the fixed-length codes in terms of bit rate but they have complex implementation. Also, they are not suitable for applications requiring fixed bit rate.

#### **2.6.3 Run Length And Huffman Coding**

After the quantization the 8\*8 matrix contains many zero coefficients particularly in the high-frequency part. These can be coded by the run-length coding efficiency. The run-length coding produce a couple of nnmbers which represent the following:

The run-length coding counts the number of zeros until a coefficient unequal zero occurs. Then one pair of numbers is generated, the first value is the count of zero coefficient and the second value is the value of the coefficient which is unequal zero

For example the run level code ( 5,20 ) represent the following 0 0 0 0 0 20.

For the run-length coding the 2-D 8\*8 matrix has to be arranged to a I-D sequence. This ordering is done by the so called Zig Zag scan way through the matrix. This Zig Zag scan way orders the quantized DCT coefficient in ascending spatial frequencies from the dc coefficient to the highest frequency coefficient. The frequency increases from the left to the right of the result matrix and also from the top to the bottom. The Zig Zag scan way is shown in Figure 1. The run-length coding is a lossless compression.

After the run-length coding each run-length couple is coded by an entropy coding. The entropy coding is also known as Huffman or variable length coding.

The appearance probability  $P_s$  of one single run-length couple is investigated. The run-length couple with the largest appearance probability is coded with a shorter bit code then run-length couple with appearance probability. The optimum code length for a symbol, Ls is given by [11]:

#### $\text{Ls} = \log_2(1/P_s)$

The entropy, which is simply the average number of bits per symbol, is given by [11] :

Entropy = 
$$
\sum P_s \log_2(1/P_s)
$$

The Huffman codes are constructed by pairing two symbols with the lowest probability combining them to a branch of a tree. The branches of the tree are assigned a1 and aO bit. The tree is developed until every symbol is covered by a branch of a tree. The Huffman codes are constructed by containing the bits of the branches, starting from the root and going back to the symbol to code.

The following example shows the results of the transformation to the frequency domain, the quantization and the run-length coding.



Quantization Results:



Results of Run-length Coding:

(0, 132) (0, -1) (0, -1) (0, 1) (0, 1) (0, -1) (I, -1) (0, -1)  $\left(0,\ 1\right)\ \ \left(2,\ 1\right)\ \ \left(7,\ -1\right)\ \ \left(13,\ -1\right)\ \ \left(1,\ 1\right)\ \ \left(10,\ 1\right)\ \ \left(0,\ 2\right)$ 

These run-length couples are coded using the Huffman coding. **In** this example a Huffman code is developed for these Run-length couples.



• Table: Probability of the run-length couples

The Huffman code can be developed with the following tree shown in figure 2.2.



Figure 2.2:

The run-length couples are coded as follows:

00111 1 1 01 01 00100 1 01 0011 00010 00000 00100 00001 00110 (50 Bit)

### **2.7 Wavelet based image coding**

There are many techniques for image coding but presently subband coding is the successful one. DCT based transform coding was popular in 1980's due to less complexity and effective bit allocation, so became the JPEG standard in image coding. JPEG image suffers from blocking artifacts. Wavelet based subband coding avoids blocking at medium bit rate, because its basis function have variable length. Wavelet based coder transforms the image into subimages using two channel filter bank.



figure 2.3: Image decomposition: first stage of the 2-D dyadic DWT

The upper left subimage is obtained by low pass filtering in both horizontal and vertical directions, represented by LL as in figure 2.3. The other subimages have high frequencies. The algorithm will assign few bits to **HH** and many bits to LL. The LL image is 4 to 5 time is iterated. These subbands are quantized and then scanned. These scanned coefficients are coded using lossless compression and then transmitted or stored. We will study in details of image compression using wavelet transform in chapter 4.

## **Chapter 3**

## **Wavelet Analysis**

In Fourier analysis frequency information can only be extracted for the complete duration of signal  $f(t)$ . Since the integral in the Fourier transform extends over all the time therefore the information rises over whole the length of signal. If there is a local oscillation at some point in  $f(t)$ , will contribute to the calculated Fourier transform  $F(w)$ . But its location in the time domain will be lost.



Figure 3.1: Discrete time wavelet series.

This disadvantage is overcome in wavelet analysis because the wavelet transform has shorter basis functions for the higher frequencies and longer basis function for the lower frequencies as shown in figure 3.1, which provides an alternative way of breaking a signal down into its constituent parts.

#### **3.1 Dilation Equations**

Now we will examine how the wavelets are generated from dilation equations. The basic function of  $\phi(x)$  is a dilated version of  $\phi(2x)$ . In dilation equation  $\phi(x)$  is expressed as a finite series of terms involving  $\phi(2x)$ . Each of these  $\phi(2x)$  terms in positioned at a different place and different argument on the horizontal axis. The basic dilution equation has a form

$$
\phi(x) = \sum_{k} c_k \phi(2x - k) \tag{3.1}
$$

where c's are numerical constants.

It is very difficult to solve equation 1 directly to find out  $\phi(x)$  therefore we have to construct  $\phi(x)$  indirectly. The simplest approach is to set up an iterative algorithm in which each new approximation is calculated by the previous one.

$$
\phi_j(x) = C_0 \phi_{j-1}(2x) + C_1 \phi_{j-1}(2x-1) + C_2 \phi_{j-1}(2x-2) + C_3 \phi_{j-1}(2x-3)
$$
(3.2)

the process of iteration is continued until  $\phi_j(x)$  and  $\phi_{j-1}(x)$  becomes almost same (indistinguishable). Consider a box function  $\phi_0 = 1 \; 0 \leq x \leq 1$ . The interval  $x = 0$  to 1 has developed a stairless function over the interval  $x = 0$  to 2. The added contribution is shown in this figure [3.2]. A particular set of co-efficient is used as defined below

$$
C_0 = (1 + \sqrt{3}) / 4 , \qquad C_1 = (3 + \sqrt{3}) / 4
$$
  
\n
$$
C_2 = (3 - \sqrt{3}) / 4 , \qquad C_3 = -(\sqrt{3} - 1) / 4 ,
$$

these co-efficient generates (see later)  $D_4$  wavelet. The D stands for Daubechies who first discovered their properties [10] .

If the iterative process is continued, the function  $\phi(x)$  approaches to limiting shape as shown in figure (3.2) and is discontinuous in nature. By magnifying the figure, we observed the graph has a fractal nature and its irregular out line remain always there. We can get smoother function by adding more terms in the dilation equation (1). This function  $\phi(x)$  is called scaling function and the corresponding wavelet function will be constructed in the next section.

The scaling function  $\phi(x)$  generated by iteration follows this matrix scheme

 $\mathbf{r}$ 

$$
[\phi_2] = \begin{bmatrix} C_0 \\ C_1 \\ C_2 & C_0 \\ C_3 & C_1 \\ C_3 & C_1 \\ C_2 & C_0 \\ C_2 & C_0 \\ C_3 & C_1 \\ C_3 & C_1 \\ C_3 & C_1 \\ C_2 & C_2 \\ C_3 & C_1 \\ C_2 & C_3 \end{bmatrix} [1] = M_2 M_1 [1]
$$



Figure 3.2: (a)  $\phi^{(1)}(t)$ .(b)  $\phi^{(2)}(t)$ .(c)  $\phi^{(3)}(t)$ .(d)  $\phi^{(4)}(t)$ . Generated by iteration

where  $M_r$  denotes a matrix of order  $(2^{r+1} + 2^r - 2) \times (2^r + 2^{r-1} - 2)$  in which each column has a submatrix of coefficients C0  $\mathrm{C}_1$  ,  $\mathrm{C}_2$  ,  $\mathrm{C}_3$  , positioned two places below of its left submatrix. As the iteration increases, the number points increases on the graph in sequence, 1, 4, 10, 22, 46,  $\cdots$   $2^{r+1} + 2$  so that after eight iteration it reaches  $2^9 + 2^8$ - 2 = 766 with each point spaced 1 /  $2^8$  = 1 / 256 unit apart. This is not most efficient method but it is simple.

### **3.2 Dilation Wavelet**

Wavelet function is derived from the corresponding scaling function with same co-efficient but in reverse order and with terms having their sign changed.

$$
\psi(x) = \sum_{k=0}^{N-1} (-1)^k c_k \phi(2x + k - N + 1)
$$
\n(3.3)

where k is positive integer and N is the total number of co-efficient. Like scaling function, it also retains the discontinuous and fractal nature but have a surprising shape for the bases function. Suppose we want to generate  $\psi(x)$  (wavelet function) from  $\phi(x)$ just after single iteration then it is generated by the following matrix scheme.

$$
\begin{bmatrix}\nC_0 \\
C_1 \\
C_2 & C_0 \\
C_3 & C_1 \\
C_2 & C_0 \\
C_3 & C_1 \\
C_2 & C_0 \\
C_3 & C_1 \\
C_3 & C_1 \\
C_3 & C_1 \\
C_2 & C_0 \\
C_3 & C_1 \\
C_2 & C_0 \\
C_3 & C_1 \\
C_3 & C_3\n\end{bmatrix} \begin{bmatrix}\n-C_3 \\
C_2 \\
-C_1 \\
C_0\n\end{bmatrix} \quad [1] = M_2M_1 [1] \quad (3.4)
$$

similarly

 $[\psi_3] = M_3 M_2 M_1$  [1]

and so on where  $M_3$  is matrix of order  $22$   $\times$   $10$ 

### **3.3 Condition for Wavelet co-efficient**

A good set of co-efficient must satisfy the following conditions.

i 
$$
\sum_{k=0}^{N-1} C_k = 2
$$
  
\nii 
$$
\sum_{k=0}^{N-1} (-1)^k k^m
$$
  
\nwhere  $m = 0, 1, 2, \dots N / 2 - 1$   
\niii 
$$
\sum_{k=0}^{N-1} C_k C_{k+2m} = 0 \qquad m \neq 0
$$
  
\niv 
$$
\sum_{k=0}^{N-1} C_k^2 = 2
$$



Figure 3.3: D4 wavelet from the scaling function

### **3.4 Discrete wavelet transforms**

The DWT algorithm was discovered by Mallat [12] and is called Mallat's pyramid algorithm or sometimes Millat' tree algorithm. We shall approach the algorithm by considering first its inverse. Suppose that the DWT has been computed to generate the sequence

$$
a=[a_0 a_1 a_2 a_3 a_4 a_5 a_6 a_7 \ldots a_{2j+k} \ldots]
$$

Suppose, for example, that we consider an expansion with the primary scaling function  $\phi(x)$  and wavelet of scale 0, 1 and 2. Then a will have  $(1 + 1 + 2 + 4) = 2<sup>3</sup>$  terms. In order to include all the wavelets at any particular scale, the total number of terms in the transform must always be a power of 2. Consider the case when there are only eight terms so that

$$
a = [a_0 \ a_1 \ a_2 \ a_3 \ a_4 \ a_5 \ a_6 \ a_7]
$$
\n
$$
(3.5)
$$

The first element  $a_0$  is the amplitude of the scaling function term  $\phi(x)$ . Since  $\phi(x)$ can be generated by iteration from a unit box over the integral  $0 \le x < 1$  (see figure),  $a_0\phi(x)$  can be generated by iteration starting from a box of height  $a_0$ . Suppose we chosen a wavelet with four coefficients. Then the first step in the iteration is from (3.4)

$$
\phi_1 = \begin{bmatrix} C_o \\ C_1 \\ C_2 \\ C_3 \end{bmatrix} [a_o]
$$

The initial box function occupied from interval  $0 \le x < 1$ , but we see in figure (3.1) that the first iteration extends over  $0 \le x < 2$ . If the part that lies outside the interval is wrapped round to fall back into the unit interval, we get

$$
\phi_1 = \begin{bmatrix} C_o + C_2 \\ C_1 + C_2 \end{bmatrix} [a_o]
$$

On taking the second iterative step, without including wrap-around, we have, from (3.4),

. Co C1 C2 Co C3 C1 Co C2 Co C<sup>1</sup> ¢ 2 = [aD] C3 C1 C<sup>2</sup> (3.6) C2 Co C3 C3 C1 C<sup>2</sup> C3

Allowing for wrap-around and adding terms at the same position in  $0 \le x < 1$ , we

can check by multiplication that this is the same as

$$
\phi_2 = \begin{bmatrix} C_a & C_2 \\ C_1 & C_3 \\ C_2 & C_0 \\ C_3 & C_1 \end{bmatrix} \begin{bmatrix} C_o + C_2 \\ C_1 + C_3 \end{bmatrix} [a_o]
$$
 (3.7)

Notice how the left-hand matrix is formed, by taking a submatrix of order  $4 \times 2$ from the left-hand matrix in  $(3.6)$  and then transposing the  $C_2$  and  $C_3$ . This recipe also applies for the matrix in  $(3.5)$  which is formed by transposing the same two elements

$$
\left[\begin{array}{c}C_0\\C_1\end{array}\right]
$$

 $\mathbf{r}$ 

For the third iterative step, the calculation is

'n

$$
\phi_3 = \begin{bmatrix}\nC_a & & & & C_2 \\
C_1 & & & & & \\
C_2 & C_o & & & & \\
C_3 & C_1 & & & & \\
C_2 & C_o & & & & \\
C_3 & C_1 & & & & \\
C_3 & C_1 & & & & \\
C_2 & C_o & & & & \\
C_3 & C_1 & & & & \\
\end{bmatrix} \begin{bmatrix}\nC_0 & C_2 \\
C_1 & C_3 \\
C_2 & C_0 \\
C_3 & C_1\n\end{bmatrix} \begin{bmatrix}\nC_0 + C_2 \\
C_1 + C_3\n\end{bmatrix} \quad [a_o]
$$
\n(3.8)

and this generates the eight ordinates in the interval  $0 \leq x < 1$  for the wrap-around scaling function. If we revise the definitions of the M matrices of the wavelet coefficients so that

$$
M_1 = \begin{bmatrix} C_o + C_2 \\ C_1 + C_3 \end{bmatrix}
$$
 of order 2 × 1 (3.9)  

$$
M_2 = \begin{bmatrix} C_o & C_2 \\ C_1 & C_3 \\ C_2 & C_0 \\ C_3 & C_1 \end{bmatrix}
$$
 of order 2<sup>2</sup> × 2 (3.10)

$$
M_3 = \begin{bmatrix} C_a & & & C_2 \\ C_1 & & & C_3 \\ C_2 & C_o & & & \\ C_3 & C_1 & & & \\ & C_2 & C_o & & \\ & & C_3 & C_1 & \\ & & & C_2 & C_o \\ & & & & C_3 & C_1 \\ & & & & & C_3 & C_1 \end{bmatrix}
$$
 (3.11)

then the algorithm for generating the contribution of  $a_0\phi(x)$  to  $f(x)$  is

$$
f^{\phi}(x) = M_3 \ M_2 \ M_1 \ a_0 \tag{3.12}
$$

or, in diagrammatic form,

$$
a_0 = f^{\phi}(1) \stackrel{M_1}{\longrightarrow} f^{\phi}(1:2) \stackrel{M_2}{\longrightarrow} f^{\phi}(1:4) \stackrel{M_3}{\longrightarrow} f^{\phi}(1:8)
$$
 (3.13)

where  $f^{\phi}(1:8)$  means an array of eight elements that represents the contribution of a<sub>0</sub> $\phi(x)$  to f(x) at  $x = 0, 1/8, 1/4, ..., 7/8$ .

Returning to the sequence  $(2)$ , consider the second term,  $a_1$ . This is the amplitude of the wavelet function  $W(x)$  which is generated from a unit box by iteration as shown in figure (3.5). The matrix operations for doing this are the same as for generating the scaling function  $\phi(x)$  except that the first step involves replacing

$$
\begin{bmatrix} C_o \\ C_1 \\ C_2 \\ C_3 \end{bmatrix} \qquad by \qquad \begin{bmatrix} -C_3 \\ C_2 \\ -C_1 \\ C_0 \end{bmatrix} \qquad (3.14)
$$

according to (3.4). The procedure for allowing for wrap-around is exactly the same as for the scaling function.

Defining

$$
G_1 = \begin{bmatrix} -C_3 - C_1 \\ C_2 + C_0 \end{bmatrix}
$$
\n(3.15)

The algorithm for generating the contribution of  $a_1 W(x)$  to  $f(x)$  is

$$
f^{(0)}(1:8) = M_3 \ M_2 \ G_1 \ a_1 \tag{3.16}
$$

or in diagrammatic form,

$$
a_1 = f^{(0)}(1) \stackrel{G_1}{\to} f^{(0)}(1:2) \stackrel{M_2}{\to} f^{(0)}(1:4) \stackrel{M_3}{\to} f^{(0)}(1:8)
$$
 (3.17)

The third term in (3.5),  $a_2$ , is the amplitude of  $W(2x)$ . Instead of the second iteration being reached by

$$
\begin{bmatrix} C_o & C_2 \ C_1 & C_3 \ C_2 & C_0 \ C_3 & C_1 \end{bmatrix} \begin{bmatrix} -C_3 - C_1 \ C_2 + C_0 \end{bmatrix} [a_2]
$$
 (3.18)

the first operation is omitted and we go straight to

$$
\begin{bmatrix}\n-C_3 \\
C_2 \\
-C_1 \\
C_0\n\end{bmatrix}\n\begin{bmatrix}\na_2\n\end{bmatrix}
$$
\n(3.19)

Thereafter the iteration proceeds as before to get

$$
f^{(1,1)}(1:8) = M_3 \begin{bmatrix} -C_3 \\ C_2 \\ -C_1 \\ C_0 \end{bmatrix} [a_2]
$$
 (3.20)

The fourth term in (3.5),  $a_3$ , is the amplitude of the translated wavelet  $W(2x - 1)$ . Allowing for wrap-around, the procedure for computing the contribution that this makes to  $f(x)$  is exactly the same as for  $a_2$  except that the elements in the first matrix are arranged in the order

$$
\begin{bmatrix}\n-C_1 \\
C_0 \\
-C_3 \\
C_2\n\end{bmatrix}
$$
\n(3.21)\n  
\n(3.22)

so that (3.20) becomes

$$
f^{(1,2)}(1:8) = M_3 \begin{bmatrix} -C_1 \\ C_0 \\ -C_3 \\ C_2 \end{bmatrix} [a_3]
$$
 (3.22)

Combining (3.20) and (3.22) then gives

$$
f^{(1,2)}(1:8) = M_3 \begin{bmatrix} -C_3 & -C_1 \ C_2 & C_0 \ -C_1 & -C_3 \ C_0 & C_2 \end{bmatrix} \begin{bmatrix} a_2 \ a_3 \end{bmatrix}
$$
 (3.23)

or putting

$$
G_2 = \begin{bmatrix} -C_3 & & \\ C_2 & & \\ -C_1 & -C_3 & \\ C_0 & C_2 & \end{bmatrix} = \begin{bmatrix} -C_3 & -C_1 \\ C_2 & C_0 \\ -C_1 & -C_3 \\ C_0 & C_2 \end{bmatrix}
$$

we get

$$
f^{(1)}(1:8) = M_3 \ G_2 \begin{bmatrix} a_2 \\ a_3 \end{bmatrix}
$$
 (3.24)

and, **in** a diagram,

$$
\begin{bmatrix} a_2 \\ a_3 \end{bmatrix} \xrightarrow{G_2} f^{(1)}(1:4) \xrightarrow{M_3} f^{(1)}(1:8) \tag{3.25}
$$

The remaining four elements of  $(3.5)$ ,  $a_4$ ,  $a_5$ ,  $a_6$ ,  $a_7$  are the amplitudes of wavelets  $\Psi(4x)$ ,  $\Psi(4x-1)$ ,  $\Psi(4x-2)$ ,  $\Psi(4x-3)$ . Each wavelet has the elements  $\begin{bmatrix} -c_3 & c_2 & -c_1 & c_0 \end{bmatrix}^t$ 

and so the single stage of calculation is

$$
f^{(2)}(1:8) = \begin{bmatrix} -C_3 & -C_1 \\ C_2 & C_0 \\ -C_1 & -C_3 \\ C_0 & C_2 \\ -C_1 & -C_3 \\ C_0 & C_2 \\ -C_1 & -C_3 \\ C_0 & C_2 \\ C_0 & C_2 \end{bmatrix} \begin{bmatrix} a_4 \\ a_5 \\ a_6 \\ a_7 \end{bmatrix}
$$
(3.26)

or, diagrammatically,

$$
\begin{bmatrix} a_4 \\ a_5 \\ a_6 \\ a_7 \end{bmatrix} \xrightarrow{G_3} f^{(2)}(1:8) \tag{3.27}
$$

combining the results we get

$$
f(1:8) = f^{(\phi)}(1:8) + f^{(0)}(1:8) + f^{(1)}(1:8) + f^{(2)}(1:8)
$$
 (3.28)

we have the final diagram below:

$$
f'(1) \stackrel{M_1}{\rightarrow} f'(1:2) \stackrel{M_2}{\rightarrow} f'(1:4) \stackrel{M_3}{\rightarrow} f'(1:8) = f
$$
\n
$$
\uparrow G_1 \uparrow G_2 \uparrow G_3 \uparrow
$$
\n
$$
a = [a(1) \quad a(2) \quad a(3:4) \quad a(5:8)]
$$
\n(3.29)

where  $a(1) = a_0$ ,  $a(2) = a_1$ ,  $a(3,4) = [a_2 \ a_3]^t$   $a(5,8) = [a_4 \ a_5 \ a_6 \ a_7]^t$ , This is the inverse of Mallat's tree algorithm [1.2]. Now consider how to break down an arbitrary function  $f(1:2^n)$  into its wavelet transform  $a(1:2^n)$ . The matrix M and G follow the orthogonality conditions, so we have

$$
\frac{1}{2} M_r^t M_r = I
$$
$$
M_r^t G_r = 0
$$
  

$$
G_r^t M_r = 0
$$
  

$$
\frac{1}{2} G_r^t G_r = I
$$

Reversing the tree and taking transpose of M and G we get the Mallat's tree algorithm.

We can decompose and reconstruct a discrete function in the same by defining a filter Land H (low pass and high pass filter).

### **3.4.1 The filter matrices L** & H

For a given vector  $f(x) = [f_1, f_2, f_3, \dots, f_n]$  where  $n = 2<sup>j</sup>$  they may be equally spaced values of a function  $f(x)$  on a unit interval. The goal is to decompose this vector at different scales. At each new level the entries cut into half. The decomposition is

$$
f = f^{\phi} + f^{(0)} + \cdot \cdot \cdot \cdot \cdot f^{j-1}
$$

the details  $f^j$  is the combination of  $2^j$  wavelets and  $f^{\phi}$  is the multiple of the scaling function  $\phi$ .

The matrix L is the fine to coarse filter and it produces a vector with half as many entries. The entries  $(L_{ij} = C_{2i-j})$  of this filter are the recursion coefficient for the scaling function. Rows 1, 2 and columns -1, 0, 1, 2 are displayed with  $N = 3$ 

$$
L = \frac{1}{2} \begin{bmatrix} C_3 & C_2 & C_1 & C_0 \\ & C_3 & C_2 & C_1 & C_0 \end{bmatrix}
$$

The beautiful thing is that the high pass filter H uses the same coefficient and in the same way this filter is associated with wavelets  $\Psi$  just as L is associated with the scaling function the filter H is defined as follow

$$
H_{ij} = (-1)^{j+1} C_{j+1-2i}
$$

for  $i = 1, 2$  and  $j = 1, 2...4$ 

$$
H = \frac{1}{2} \left[ \begin{array}{cccc} C_0 & -C_1 & C_2 & -C_3 \\ & C_0 & -C_1 & C_2 & -C_3 \end{array} \right]
$$

These indices are used to match the Haar different are also possible. The important points are

$$
H \quad L^* = 0
$$
  

$$
L \quad L^* = I \qquad \& \quad H \quad H^* = I
$$
  
and 
$$
L^* \quad L + H^* \quad H = I
$$

## **3.4.2 Decomposition**

The decomposition of a discrete function f in to Mallat's Pyramid algorithm is as follow starting from

$$
a^{j} = f \qquad for \quad i = J, \dots, 1
$$
  

$$
a^{j-1} = La^{j} \qquad and \qquad b^{j-1} = Ha^{j}
$$

the full decomposition is represented by a tree of filters

$$
a^{J} \xrightarrow{L} a^{J-1} \xrightarrow{L} a^{J-2} \cdots \xrightarrow{L} a^{0}
$$
  

$$
\searrow b^{J-1} \searrow b^{J-2} \cdots \searrow b^{0}
$$

### **3.4.3 Reconstruction**

The decomposition of discrete function at different scale can be recovered by starting from  $a^0$  and  $b^0, \cdots, b^{J-1}$  for  $j = 1, \ldots, J$ 

$$
a^j = La^{j-1} + H^*b^{j-1}
$$

The reconstruction goes from the branches of the tree back to the root

$$
a^{0} \xrightarrow{L^{*}} a_{1} \xrightarrow{L^{*}} a^{2} \cdots \xrightarrow{L^{*}} a^{J} = f
$$

$$
b^{0} \nearrow_{H^{*}} b_{1} \nearrow_{H^{*}} \cdots \nearrow_{H^{*}}
$$

## **3.5 Wavelet and Multiresolution Analysis**

Multiresolution will be described first for subspaces  $V_j$  and  $W_j$ . The scaling spaces  $V_j$  are increasing. The wavelet space  $W_j$  is the difference between  $V_j$  and  $V_{j+1}$ . The sum of  $V_j$ and  $W_j$  is  $V_{j+1}$ . Then these extra conditions involving dilation to  $2<sub>t</sub>$  and translation to  $t - k$  define a genuine multiresolution:

If  $f(t)$  is in  $V_j$  then  $f(t)$  and  $f(2t)$  and all  $f(t-k)$  and  $f(2t-k)$  are in  $V_{j+1}$ .

In the end, one wavelet generates a whole basis. The functions  $w(2^j t - k)$  come by dilation and translation (all  $j$  and all  $k$ ). There are six steps toward this goal, and we take them one at a time:

- 1. An increasing sequence of subspaces  $V_j$  (complete in  $L^2$ ).
- 2. The wavelet subspace  $W_j$  that gives  $V_j + W_j = V_{j+1}$
- 3. The dilation requirement from  $f(t)$  in  $V_j$  to  $f(2t)$  in  $V_{j+1}$
- 4. The basis  $\phi(t k)$  for  $V_0$  and  $w(t k)$  for  $W_0$
- 5. The basis  $\phi(2^{j}t k)$  for  $V_j$  and  $w(2^{j}t k)$  for  $W_j$
- 6. The basis of all wavelets  $w(2^{j}t k)$  for the whole space  $L^{2}$

### **3.5.1 A Scale Of Subspaces**

Each  $V_j$  is contained in the next subspace  $V_{j+1}$ . A function in one subspace is in all the higher (finer) subspaces:

$$
V_0 \subset V_1 \ldots \subset V_j \subset V_{j+1} \subset \ldots
$$

A function  $f(t)$  in the whole space has a piece in each subspace. Those pieces contain more and more of the full information in  $f(t)$ . The piece in  $V_j$  is  $f_j(t)$ . One requirement on the sequence of subspaces is completeness.

$$
f_j(t) \to f(t)
$$
 as  $j \to \infty$ 

### **3.5.2 The Dilation Requirernent**

So for we have an increasing and complete scalar of spaces. Each is  $V_j$  contained in the next  $V_{j+1}$ . For multiresolution, the crucial word *scale* carries an additional meaning.  $V_{j+1}$  consists of all re scaled functions in  $V_{j+1}$ 

*Dilation* :  $f(t)$  is in  $V_j \iff f(2t)$  is in  $V_{j+1}$ .

In addition to completeness as  $j \to \infty$ , we require emptiness as  $j \to -\infty$ :

 $\cap V_j = \{0\}$  and  $\cup V_j =$  whole space.

Emptiness means that  $|| f_j(t) || \to 0$  as  $j \to -\infty$ . Completeness means that  $f_j(t) \to$  $f(t)$  as  $j \rightarrow \infty$ . The detail:

 $\Delta f_j = f_{j+1} - f_j$  belongs to  $W_j$  and we still have

 $V_i \otimes W_j = V_{j+1}$ 

This can be orthogonal sum, with  $\Delta f_j$  orthogonal to  $f_j$ . It must be a direct sum, with  $V_j \cap W_j = \{0\}$ . The construction of  $f(t)$  from its details  $\Delta f_j$  can start at  $j = 0$  as before, or it can start at  $j = -\infty$ :

### **3.5.3 The Translation Requirement And The Basis**

Instead of rescaling  $f(t)$ , we now shift its graph. This is *translation*, and it leads to the fundamental requirement of time-invariance in signal processing. The subspaces are *shift* - *invariant:* 

*If f<sub>j</sub>*(*t*) *is in*  $V_j$  *then so are its translates f<sub>j</sub>*(*t* - *k*).

Suppose  $f(t)$  is in  $V_0$ . Then  $f(2t)$  is in  $V_1$  and so is  $f(2t - k)$ . By induction,  $f(2<sup>j</sup>t)$  is in  $V_j$  and so is  $f(2^jt - k)$ . Dilation and translation are now built in.

## **Chapter 4**

# **Image Compression Using Wavelet Transform**

The main obstacle for many application of digital images, i.e,acquisition, data storage, printing, and display, is the huge amount of data required to represent an image directly. Such an image needs to be compressed for storage or transmissions. The actual compression ratio can vary from 100:1 to 2:1 depending on the specific application and encoder/decoder complexity. State-of-the art techniques can compressed typical images by a factor of 10 to 50 without significantly effecting the image quality, depending on the technique applied. There are many techniques for image coding, we use the most successful, i.e,wavelet based subband image coding that avoid blocking artifact (which is the main disadvantage of JPEG standard) at medium bit rate, because the variable length of its basis functions.

## **4.1 Subband Coding**

Subband coding (SBC) is another form of frequency decomposition. In SBC, a signal is decomposed into a number of equal- or unequal-frequency bands using filter banks that have been developed recently  $[13, 14, 15]$ . In fact by using perfect reconstruction filter

banks [5], the original signal going through the frequency decomposition-subsamplinginterpolation-synthesis process can be fully recovered. The philosophy behind SBC is that coding techniques compatible with the frequency bands can be applied. signal components in high-frequency bands can be either dropped out or coarsely quantized.

Subband coding is also ideally suited for progressive image transmission (PIT) as bits related to lower bands can be transmitted first, followed by those related to the upper bands.

## **4.2 Subband Image Coding**

A popular approach to subband image coding [16] is to map the image into four equal subbands in the 2D frequency domain as shown in figure 4.1. In general, a transform (such as DCT) is applied to the lowest subband, followed by quantization and VLC. The remaining subbands are coarsely quantized.



Figure 4.1: Mapping of an image into four equal sbbands in the 2-D frequency domain (L : low frequency, H: high frequency).



Figure 4.2: Unequal subband decomposition proposed by Bellcore for the ATM / SCONET / H-4-ISDN HDTV project **[17J** 

## **4.3 DWT And Subband Coding**

The DWT is known to be generated by a cascade of filter banks and the DWT is essentially the well-known subband decomposition [18]. The advantages of the DWT comes from the trade-off between spatial and frequency resolution, as the DWT has shorter basis functions for the higher frequencies as shown in the fig. Therefore, DWT has a versatile time-frequency localization due to a pyramid like multiresolution decomposition.

The 2-D wavelet transform is implemented independently, first in the horizontal direction and then in the vertical direction [12]. This method of applying the DWT to a two dimensional signal that is sampled on a rectangular lattice is called *the dyadic*  2-D DWT. The dyadic 2-D DWT decomposes the original image into four subbands, as shown in figure  $4.3(a)$ .



Figure 4.3: Image decomposition (a) first stage of the 2-D dyadic DWT (b) third stage of the 2-D dyadic DWT. V and W mean subspaces and difference subspaces, respectively.

One of the subband contains all the low-pass information (LL) and another all the high-pass information (HH). The other two subbands are a horizontal high-pass band containing vertical low-pass information (LH) and a vertical high-pass containing horizontal low-pass information (HL). For a full DWT the decomposition is repeated several times on the LL part [19]. Fig.4.3(b) shows the decomposition of the image after three iterations. There are a total of ten sub-images, which are used to encode the original image.  $V$  and  $W$  were defined as subspaces and spaces, respectively. This decomposition provides sub-images corresponding to different resolution levels and orientations.

Figure 4.4 shows the analysis section of a two-dimensional (2D) separable filter bank, where the first image rows are passed through the 2-channel filter bank and then, the columns are processed. The analysis section can be viewed as a  $2 \times 2$  transform applied to the image (note that each subband has one fourth of the samples in the original signal). Also, the synthesis section can be viewed as a  $2\times 2$  inverse transform. Note that only the low-pass subband is connected to next stage transform. The inverse transform is, of course, accomplished by reversing the paths and the transforms.

To reconstruct the image from the subbands, one may choose the set of low-pass and high-pass filters shown in figure 4.4, so as to provide perfect reconstruction.



Figure 4.4: One step (first level) of a 2-D wavelet decomposition.

### **4.3.1 Entropy Coding:**

After bit allocation and quantization, we have subimages with discrete levels represented by integers. How do we store or transmit these subimages? Many high pass coefficients are zero after quantization. These coefficients should be grouped so that the entropy coder can take full advantage of long strings of zeros. This is accomplished by *scanning.* 

Run-length coding or Huffman coding or a combination should be used to reduce the redundancy of the images. We will discuss the baseline entropy coder which is a combination. JPEG also uses the baseline coding method.

### **4.3.2 Scanning Of The D iscrete Wavelet Coefficients:**

To demonstrate scanning, consider the three level wavelet transform in Figure 4.5. Subband 2,5 and 8 are highly correlated since 2 is the coarse approximation of 5 and 5 is the coarse approximation of 8. Suppose the pixel at the upper left corner of subband 2 is zero. Then it is very likely that the pixels in a  $2 \times 2$  shaded square of subband 5 are *zero.* Similarly, the pixels in a  $4 \times 4$  shaded square of subband 8 are probably zero. One can these group pixels into an "AC sequence" of length  $21(=1+4+16)$  by vertically scanning the shaded squares. Figure also shows the scanning patterns of subbands 3, 6 and 9 (horizontal) and for subbands 4, 7 and 10 (diagonal). When the original image has size 32, the 16 pixels in each subband2, 3, and 4 give 48 AC sequences. The low frequency band is scanned horizontally and grouped into the DC sequence of length 16. This scanning method is similar to the *zero-tree coder* proposed by [Shapiro].



Figure 4.5: Scaning method used in the discrete wavelet decomposition.

### 4.3.3 Sequential Baseline Coding:

After scanning the quantized subimages, we have a set of DC and AC sequences to be stored. The baseline coder takes advantage of the correlation in the AC sequences. This algorithm combines Sequential Baseline Coding and Huffman entropy coding. The basic principle is similar to the JPEG coder, but does not restrict to the DCT.

### *Coding Of The AC Sequence:*

The sequence  $900201000 - 30030 - 1 - 1$  has strings of zeros interlacing with nonzero's. An efficient representation remembers the number of zeros before each nonzero. *Symbol-l* is the pair *(runlength, size)* where *runlength* specifies the number of preceding zeros. *Size* determines the number of bits to encode the current nonzero. *Symbol-2* gives the *(amplitude)* of the nonzero: *Size = n* corresponds to *amplitude* less than  $2^n$ (but not less than  $2^{n-1}$ ).

This representation of the example gives a string of *Symbol-i* and Symbol-2:

 $(0,1)1 (2,2)2 (1,1)1 (3,2)-3 (2,2)3 (1,1)-1 (0,1)-1 (0,0)$ :

Note the terminal *symbol-1* (0, 0) at the end. Also (1, 1) and (0, 1) and (2, 2) occur twice in the string. Huffman entropy coding as shown in figure 4.6 can exploit this redundancy in *symbol-i.* 





### *Coding OJ The DC Sequence:*

DC coefficients measure the average energy of the input signal. There are usually a strong correlation between neighbouring coefficients. For efficiency, we use *differential coding:*  save the first coefficient and then the *differences between successive coefficients* . These are coded as for AC coefficients. Since one would not expect long zero strings, *runlength*  is not used. *Symbol-i* only gives *size.* 

Three error measures are often used to compare coders and perceptual quality:

Mean Square Error  $MSE = \frac{1}{mn} \sum_{m=0}^{m-1} \sum_{n=0}^{n-1} | x(m,n) - \tilde{x}(m,n) |^2$ Peak Signal Noise Ratio  $PSNR = 10 \log_{10}(\frac{255^2}{MSE})$ 

Lena	<b>PSNR</b>	$32.6$ 33.5 34.1 35.5 36.4 37.0 37.8 38.5 39.0 39.6				
Barbara PSNR		25.8   26.9   27.8   29.1   29.9   31.6   32.1   33.1   33.7   34.8				
Goldhill PSNR		$29.8$ 30.4 31.0 32.0 32.9 33.8 34.1 34.9 35.6 36.1				

Table 4.1: PSNR and Bit rate of Lena, Barbara and Goldhill

Maximum Error *MaxError* =Max $|x(m, n) - \widehat{x}(m, n)|$ 

The image is  $M\times N$  . The MSE and PSNR are directly related, and one normally uses PSNR to measure the coder's objective performance. At high rate, images with PSNR above 32 db are considered to be perceptually lossless [5]. At medium and low rates, the PSNR does not agree with the quality of the image.



Figure 4.10: Comparisons of wavelet based image coding and JPEG, in bitrate and PSNR for the 512 x 512 Lena image.

×



Figure 4.11: Comparisons of wavelet based image coding, in bitrate and PSNR for the 512 x 512 Lena, Barbara and Goldhill images.



Figure 4.12: Comparisons of wavelet based image coding, in bitrate, RMS error, for the 512 x 512, Barbara, Lena, and Goldhill images

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# *Appendix*



#### encode.cpp

//ENCODE.com #include <stdio.h> #include <stdarg.h> #include <stdlib.h> #include <math.h> #include <assert.h> #include <iostream.h> #include <fstream.h> \*include "trans.hh" #include "coeffset.hh" #include "allocr.hh" #include "quant.hh" void compress (Image \*image, Wavelet \*wavelet, int nStages, int capacity, Real p, Real \*weight, int paramPrecision, int budget, int nQuant, Real minStepSize, char \*filename, int monolaver) : int main void) char input\_name[50], output\_name[50]; Real Ratio: #ifdef PGM printf("Please enter the Input file name\n"),  $scanf(*ks", kinnut, nam)$ : printf("Please enter the Output file name\n"), scanf ("%s", &output name) : printf("Please enter the Compression Ratio\n"). scanf ("%f", &Ratio); char "infile name = input name: char \*outfile name = output name: Real ratio = Ratio: // Load the image to be coded Image \*image = new Image (infile\_name); int budget = (int)((Real)(image->hsize\*image->vsize)/ratio); // (assumes & bit pix  $e(s)$ printf ("Reading %d x %d image %s, writing %s\nCompression ratio %g:1\n", image->hsize, image->vsize, infile name, outfile name, ratio); telse if  $(\text{area} \leq 6)$  ( fprintf (stderr, \*Usage: %s [image] [width] [height] [output] [ratio] \n\*, program); forintf (stderr, "image: image to be compressed (in RAW format) \n") : forintf (stderr, "width, height: width and height of image to be compressed\n"); fprintf (stderr, "output: name of compressed image\n"); fprintf (stderr,

"ratio: compression ratio\n");  $ext(0);$ char \*infile name =  $aravl1$ :  $int \; hist$  = atoi(argv(2));  $int vsize = atoi (arav[3]):$ char \*outfile name =  $aravl$ Real ratio =  $atof(aray[5]):$ int budget = (int)({Real)(hsize\*vsize)/ratio); // (assumes 8 bit pixels) printf ("Reading %d x %d image %s, writing %s\nCompression ratio %g:1\n", hsize, vsize, infile name, outfile name, ratio); // Load the image to be coded Image \*image = new Image (infile\_name, hsize, vsize); tendif // Create a new wavelet from the 7/9 Antonini filterset Wavelet "wavelet = new Wavelet (&Antonini): // Coding parameters Real  $p = 2.0$ : // exponent for L^p error metric int  $ns$ tages = 5: // # of stages in the wavelet transform int capacity =  $512:$ // capacity of histogram for arithmetic coder int paramPrecision =  $4$ ; // precision for stored quantizer parameters // # of quantizers to examine for allocation int nouant =  $10$ ; Real minStepSize = 0.05; // smallest quantizer step to consider int monolayer =  $FALSE$ ; // TRUE for non-embedded uniform quantizer // FALSE for multilayer quantizer Real \*weight = // perceptual weights for coefficient sets new Real(3\*nStages+1): // for now give all sets equal weight for (int  $i = 0$ ;  $i < 3$ \*nStages+1;  $i \rightarrow i$  $W \approx 1 \pi h t$ ,  $111 = 1.0$ . compress (image, wavelet, nStages, capacity, p, weight, paramPrecision, budget, nQuant, minStepSize, outfile\_name, monolaver): delete () weight: delete image: delete wavelet: return 0: void combress (Image \*image, Wavelet \*wavelet, int nStages, int capacity, Real p, Real \*weight, int paramPrecision, int budget, int nouant, Real minStepSize, char \*filename, int monolayer) int i: // Compute the wavelet transform of the given image WaveletTransform \*transform = new WaveletTransform (wavelet, image, nStages); // For each subband allocate a CoeffSet, an error metric, an // EntropyCoder, and a Quantizer int nSets = transform->nSubbands: CoeffSet \*\* coeff = new CoeffSet\* InSets1: ErrorMetric \*\*err = new ErrorMetric\* InSetsl:

EntropyCoder \*\*entropy = new EntropyCoder\* (nSets);

### encode.cpp

**Quantizer w·quant new Quantizer\* InSets);**   $for (i = 0; i < nSets; i++)$  (  $err[i] = new LDFror (p)$ ; if  $(monolaver)$   $l$ J *II* **Partition the wavelet transformed coefficients into subbands --** *II* **Use uniform quantizer and single layer escape coder for each**  *I I* subband entropy[i] = new EscapeCoder (capacity); *II* Assume all subbands have pdf's centered around 0 except the *II* low pass subband 0 quant[i] = new UniformQuant ((MonoLayerCoder \*)entropy[i], else ( **paramPrecision, i != 0, err(i]);**  *II* **Use a layered quantizer with dead zones and a layered entropy**  *' II* coder for each subband entropy  $[i]$  = new LayerCoder (nQuant, i  $!= 0$ , capacity); I/ Assume all subbands have pdf's centered around 0 except the *II* low pass subband 0 quant  $[i]$  = new LaverQuant ((MultiLaverCoder \*)entropy $[i]$ , **paramPrecision, i != 0, nQuant, err{i]);**  *II* **each subband will have a different quantizer**   $coeff[i]$  = new CoeffSet (transform->subband(i), transform->subbandSizeIil, quant[i]); *II For* **each** sub band **determine the rate and distortion for each of**  *II* **the possible nQuanc quantizers**  coeff[i] ->qetRateDist (nQuant, minStepSize); **Allocator \*allocator = new Allocator ();**  *II* **Use rate <sup>l</sup> distortion information for each subband to find bit**  *II* **allocation that minimizes total (weighted) distortion subject**  II to a byte budget **budget -= nSets w 4;** *II* **subtract off approximate size of header info**  allocator->optimalAllocate (coeff, nSets, budget, TRUE, weight); printf ("Target rate = %d bytes\n", budget); // Display the resulting allocation **allocator - >print (coefE, nSets) ;**  // Open output file ofstream outfile (filename, ios::out | ios::trunc | ios::binary); if (!outfile) ( *error* **( · Unable to open file %s·, filename);**  *II* **Create** *1/0* **interface object for arithmetic coder Encoder wencoder = new Encoder (outfile);**  *II* **Write image** *size* **to output file encoder->writePositive (image - >hsize); encoder->writePositive (image->vsize);**  *II* **printf ( ·hsize = %d, vsize = %d\n-, image->hsize , image->vsize);**   $for$   $(i = 0; j < nSets; i++)$  ( // Write quantizer parameters for each subband to file  $\cosh\left(\frac{1}{2}\right) - \text{wrlteHeader}$  (encoder, allocator->precision(i));

*II* **Quantize and write entropy coded coefficients for each subband**  coeff[il->encode (encoder, allocator->precision[i]);

*II* Flush bits from arithmetic coder and close file encoder->flush [I; **delet e encoder;**  outfile,close ();

*1/* Clean up for  $(i = 0; i < nSets; i++)$  ( delete err[i]; delete entropy[i]; delete quant[i]; delete coeff[i]; delete [J err; delete [I entropy; delete [J quant; **delete** [J coef f; **delete allocatorj delete transformj** 



### decode.cpp

**#include <stdio .h>**  #include <stdarg . h> .include <stdlib.h>  $i$ include  $\epsilon$ math  $h$ **#include <assert . h>**   $*$ include <iostream.h> **#include <fstream.h> #include - trans .hh · #include · coeffset . hh <sup>M</sup>** 'include "allocr.hh" 'include "quant.hh " void decompress (Image \*\*image, Wavelet \*wavelet, int nStages, int capacity, int paramPrecision, char \*filename, **int nQuant , int monolayer)** *<sup>i</sup>* **int main**  void) ( char input file[50], output\_file[50];  $print('Please Enter The Code of File\n$  ;  $scan f$  ( $ks$ ,  $\delta$ input file) ; printf ( "Please Enter The Output File\n") ; scanf ("%s", &output\_file); **char** ~infile\_name **<sup>=</sup>input\_fil e ;**  char \*outfile\_name = output\_file; **printf (-Reading compressed image %s, writing %s\ n\n ", infile\_name, outfile\_ nama) ;**  *II* **Create a new wavelet from the** *7/9* **Antonini fil te rset**   $W$ <sup>avelet</sup> \*wavelet =  $n$ ew  $W$ avelet (&Antonini); **1/ Coding parameters /1 # of stages in the wavelet transform i nt nStages = 5 ; 1/ capacity o f histogram for arithmetic coder int capacity = 512; int paramPrecision = 4' 1/ precision for stored quantizer parameters 1/ # of quantizers to examine for allocation int nQuant = 10;**  *II* **TRUE for non - embedded uniform quan ti zer i nt monolayer = FALSE;**  */1* **FALS2 for multilayer quantizer Image \*reconstruct; decompress {&reconstruct, wavelet. nSta ges, capacity, pa ramPrecision, infile\_name, nQuant, monolayer} ;**  <sup>i</sup>ifdef PGM  $reconstruct->savePGM (outfile_name);$ **#else reconstruct->saveRaw** (out file\_name) ; **#endif delete reconstruct; delet.e wavelet; return 0;** 

```
void decompress (Image **image, Wa ve let ~wavelet , i nt nStages, 
                 int capacity, int paramPrecision, char *filename, 
                 int maxQuant, int monolayer) 
 int i; 
  II open compressed image file 
 ifstream infile (filename, ios::in | ios::nocreate | ios::binary);
 if ( !infile) ( 
   error ( MUnable to open file %s ·, filename ): 
  1/ Create 1/0 interface object for a r ithmetic decoder 
  Decoder *decoder = new Decoder (infile); 
  II Read i mage dime nsions from file 
  int hsize = decoder->readPositive ( )i 
 int \text{vsize } \pm \text{decoder} > r eadPositive () :
  II printf ("hsize = \delta d, vsize = \delta d \nabla", hsize, vsize);
 VJaveletTransform * transforrn. = 
   new WaveletTransform (wavelet , hsize, vsi ze , nStages); 
  // For each subband allocate a CoeffSet, an En t ropyCoder, and a 
  II Quanti z er (don't need to know anything about errors here) 
  int nSets = transform->nSubbands ; 
  CoeffSet *""coeff = new CoeffSet* (nSetsl; 
  EntropyCode ** entropy = new EntropyCoder* (nSets);
  Quantizer **quant. = ne\"! QIJantizer* (nSets]; 
  II Quantizer precision for each subband 
  int *precision = new int [nSets];
  for (i = 0; i < nSets; i++ ) { 
    if (mono laye r) ( 
      II Use uniform quanti zer and single layer escape coder for each 
      II subband 
      entropy[i] = new EscapeCoder (capacity);
      1/ ASSllme all subbands have pdf's cen t ered around 0 except the 
      II low pass subband 0 
      quant(i] = new UniformQuant. 
        ((Mono LayerCoder *)ent.ropy(ij, pa ramPrec isi on, i ! =0); 
    else ( 
      en tropy lil = new LayerCoder (maxQuant , i != 0, capacity) ; 
      11 Assume all subbands have pdf's centered around 0 except the
      II low pass subband 0 
      quant [i] = new LayerQuant
        ((MultiLayerCoder * )entropy[i j, paramPrec i sion. i != 0, maxQ1lant); 
     } 
    // Indicate that each set of coefficients to be read corresponds
    II to a subband
    coeff[i] = new CoeffSet (transform->subband(i),
                               transform->subbandSize[i], quant[i]);for (i = 0; i < nSets; i++)// Read quantizer parameters for each subband 
    coeffliJ->readHeader (decoder, precision[i1);
```
for  $(i = 0; i < nSets; i++)$  {

## .decode.cpp

**II Read, decode, and dequantize coefficients for each subband**  coeff[i]->decode (decoder, precision[i]);  $\mathbf{V}$ 

```
// Close file
del ete decoder; 
infile.close ();
```
// Clean up for  $(i = 0; i < nSets; i++)$  { delete entropy[i); delete quant[i]; delete coeff[i]; **delete** [) entropy; delete [] quant; delete [] coeff; delete [) **precision;** 

**II Allocate image and invert transform -image;;;; new Image (hsize, vsize); transEorm->invert (\*image); delete transform;** 

### **a11ocato.cpp**

I' Allocato . cpp With the same of the company of the consideration of the company of the company of the company **#include <stdio.h> #include <iostream.h>**  iinclude <math . h> iinclude "global . hh" iinclude "entropy .hh " iinclude "quant . hh " **#include ·coder.hh-** !include "a llocr.hh " Allocator : - Allocator () {  $precision = NUL.$  $\overline{f}$  , and a set of the contract of the c Allocator: - "Allocator () ( if (precision != NULL) **delete (J precision ; void Allocator : :resetPrecision (inc nSets)**  ( **if (precision !;;: NULL) delete [J precision ; precision;;: new int [nSets1; for (inc i ;;: 0; i < nSecs; i++J**   $presion(1) = 0$ ; **void Allocator: :optlrnalAllocate (CoeffSet** ~~coeff, **int nSets, int budget, int augment, Real "'weight)**   $Real bltBudget = 8*budget;$ Real lambda, lambdaLow, lambdaHigh; Real rateLow, rateHigh, currentRate; Real distLow, distHigh, currentDist; **resetPrecision (nSets);**  lambdaLow = 0.0; allocateLambda (coeff, nSets, lambdaLow, rateLow, distLow, weight); if (rateLow < bitBudget) // this uses the largest possible # of bits return:  $\frac{1}{2}$  -- if this is within the budget, do it **return; 1/ -- if this is within the budget, do it**  lambdaHigh = 1000000.0; Real lastRateHigh = -1; do ( **II try to use the smallest possible" of bits**  allocateLambda (coeff, nSets, lambdaHigh, rateHigh, distHigh, weight);

```
II if this is still> bitBudget, try again wI larger lambda 
   if (rateHigh > bitBudget && lastRateHigh != rateHigh) ( 
     lambda Low = lambdaHigh; 
     rateLow = rateHigh; 
     distLow = distHigh; 
     lambdaHigh *= 10.0; 
 ) while (rateHigh > bitBudget && lastRateHigh != rateHigh) :
  II give up when changing lambda has no effect on things 
 if (lastRateHigh == rateHigh)
   return; 
  II Note rateLow will be> rateHigh 
 if (rateLow < bitBudget) 
    error ('Failed to bracket bit budget = %d: rateLow = %g rateHigh = %g\n",
          budget, rateLow, rateHigh) ;
 while (lambdaHigh - lambdalow > 0.01) (
   lambda = (lambdaLow + lambdaHigh) 12 .0; 
   allocateLambda (coeff, nSets, lambda, currentRate, currentDist, weight); 
   if (c1.lrrentRate> bitBudget) 
     lambda Low = lambda; 
   else 
      lambdaHigh = lambda; 
 if (currentRate > bitBudget) 
   lambda = lambdaHigh; 
   allocateLambda (coeff, nSets, lambda, currentRate, currentDist, weight);
 if (augment) 
    qreedyAugment (coeff, nSets, bitBudget-currentRate, weight) :
void Allocator: :allocateLambda (CoeffSet **coeff, inc nSets, 
                                Real lambda, Real &opcimalRate, 
                                Real &optimalDist, Real *weight)
   inc i, j; 
  Real G, minG, minRate, minDist;
  optimalRate = optimalDist = 0.0;11 want to minimize G = distortion + lambda * rate
   II loop through all rate-distortion curves 
   for (i = 0; i < nSets; i++)minG = minRate = minDist = MaxReal;for (j = 0; j < coeff(i)->nQuant; j_{j+1}------------*1 
      G = weight[i]*coeff[i]->dist[j] + lambda * coeff[i]->rate[j];
      if (G < minG) (
        minG = G; 
        minRate = coeff{iJ->rate(j]; 
        minDist = weight[i]*coeff[i]-addist[j];precision[i] = j;
```


### allocato.cpp

```
optimalRate += minRate:
    optimalDist += minDist;
void Allocator::greedyAugment (CoeffSet **coeff, int nSets,
                           Real bitsLeft, Real *weight)
  int bestSet, newPrecision = -1:
 Real delta, maxDelta, bestDeltaDist, bestDeltaRate = 0:
  do {
   bestSet = -1:
   maxDelta = 0:
   // Find best coeff set to augment
   for (int i = 0; i < nSets; i++) {
     for (int j = precision[i]+1; j < coeff[i]->nQuant; j++) (
       Real deltaRate = coeff[i]->rate[j] -
         coefflil->rate[precision[i]]:
       Real deltaDist = - weight[i]*(coeff[i]->dist[i] -
         coeff[i]->distiprecision[i]]);
       if (deltaRate != 0 && deltaRate <= bitsLeft) (
         delta = deltaDist / deltaRate:
         if (delta > maxDelta) (
          maxDaIta = delta:
          hestDeltaRate = deltaRate:
          bestDeltaDist = deltaDist:
          bestSet = i:
          newPrecision = j;
   if (bestSet != -1) {
     precision [bestSet] = newPrecision;
     bitsLeft -= bestDeltaRate;
 ) while (bestSet != -1):
void Allocator::print (CoeffSet **coeff, int nSets)
 Real totalRate = 0, totalDist = 0;
 int totalData = 0.
  for (int i = 0; i \lt pSets; i \leftrightarrow j (
  totalRate += coeff[i]->rate[precision[i]];
  totalDist += coeff[i]->dist[precision[i]];
  totalData += coeff(i)->nData;
 Real rms = sqrt(totalDist/(Real)totalData);
```

```
Real psnr = 20.0 * log(255.0/\text{rms})/log(10.0):
```
#### printf ("\n"):

// printf ("total rate =  $% g(n)$ ", totalRate/8.0): // printf ("total dist =  $\sqrt[3]{n}$ ", totalDist); //printf ("total coeffs = %d\n", totalData); printf ("RMS error =  $\alpha \nvert n$ ", rms): printf (\*PSNR (transform domain) = %g\n', psnr);

```
11#include <iostream.h>
#include <iomanip.h>
#include <assert.h>
#include <math.h>
#include "global.hh"
*include "BitIO.h"
#include "iHisto.h"
#include "Arith.h"
const int CodingValues:: CodeValueBits = 16:
const long CodingValues:: MaxFreq = ((long)1 << (CodeValueBits - 2)) - 1:
const long CodingValues:: One = ((\text{long})1 \ll \text{CodeValueBits}) - 1:
const long CodingValues: : Otr = One / 4 + 1;
const long CodingValues: : Half = 2 * Qtr;
const long CodingValues:: ThreeOtr = 3 * Otr:
ArithEncoder:: ArithEncoder(BitOut &bo) : output(bo)
    1ow = 0hich = OnebitsToFollow = 0:piov
ArithEncoder::flush(void)
    for (int i = 0; i < CodeValueBits; i \leftrightarrow jIf (1ow \geq Half)bpf(1);
            1ow = HalfJ else (
            bpf(0);1ow = 2 * 1owoutput.flush();
void ArithEncoder::Encode(int count, int countLeft, int countTot)
         cerr << "Encode(" << count << ", " << countLeft << ", " << countTot << ")\n
    11A.
```
assert(count); long range = high - low - 1; high =  $low + (range * (countLeft + count)) / countTot - 1;$  $low = low + (range * countLeft) / countTot;$ while  $(1)$   $($ if (high  $<$  Half) {  $bpf(0)$ : ) else if (low >= Half) ( bpf $(1)$ ;  $low -= Hallt$  $high = Half$ ) else if (low >= Otr && high < ThreeOtr) ( bitsToFollow++:

 $1ow = OCF:$ 

```
arith.cpp
                          high - 0tr:
                      \int also \intbreak.
                      1ow = 2 * 1ow:high = 2 * high + 1.
              ArithDecoder::ArithDecoder(BitIn &bi) : input(bi)
                  1ow = 0:
                  hich = One:value = 0:
                  for (int i = 0; i < CodeValueBits; i++) {
                      value = \{value \ll 1\} + input. input bit():
              int
              ArithDecoder::Decode(iHistogram &h)
                  long range:
                  int cum:
                  int answer:
                  int ct. ctLeft. ctTotal:
                  ctTotal = h.TotalCount();range = high - 1ow + 1:
                  cum = ((\text{long}) (\text{value - low}) + 1) * ctTotal - 1) / range;
                  answer = h.Symbol(cum):
                  et = h.Count(answer); ctLeft = h.LeftCount(answer);
                  // cerr << "Decoder : cum= " << cum << "-> " << answer << "\n";
                  11 cerr \ll \approxct = ' << ct << ' ctLeft = ' << ctLeft << ' ctTotal = ' <<
               ctTotal << "Vn":high = low + (range * (ctLeft + ct)) / ctTotal - 1;
                  low = low + (range * ctleft) / ctfotal;while (1) (if (high \langle Half) \langle\frac{1}{2} else if (low >= Half) (
                       value = Ha1f:low = <code>Half+</code>high -= Half:
                      else if (1ow >= 0tr&& high < ThreeQtr) {
                          value = 0trtlow == 0tr:high -= Otr:
                      else break;
                      1ow = 2 * 1owhigh = 2 * high + 1;
                      value = 2 * value + input import_bit();return answer:
```
### coder.cpp

//coder.cpp #include <iostream.h> #include <stdio.h> #include <math.h> #include "global.hh" #include "coder.hh" Encoder:: Encoder (ostream &out, ostream &log)  $\mathcal{I}$ bitout = new BitOut (out,  $log$ );  $arith = new Arithmetic()$ ; inteoder = new CdeltaEncode (bitout) :  $\rightarrow$ Encoder: "Encoder () delete intcoder: delete arith: delete bitout: Decoder::Decoder (istream &in, ostream &log)  $\tau$ bitin = new BitIn (in, log); intcoder = new CdeltaDecode (bitin):  $arith = NULL$  $\mathcal{N}$ Decoder: : "Decoder () delete introder: if (arith != NULL) delete arith: delete bitin:  filter.cpp

// Filter.cpp **linclude <stdio.h>**  tinclude <stdarg.h> linclude <stdlib . h> **linclude <math . h> 'include <assert.h>**  linclude "wavelet.hh" Real HaarCoeffs [[ *{l.0/Sqrt2, 1 . 0/Sqrt2* }; *II* A few Daubechies filters Real Daub4Coeffs [] = [ 0.4829629131445341, 0.8365163037378077, 0.2241438680420134, - 0.1294095225512603 } ; Real Daub6Coeffs [] = ( 0.3326705529500825, 0.8068915093110924, 0.4598775021184914, -0.1350110200102546, -0.0854412738820267, 0.0352262918857095}; Real Daub8Coeffs [ ] = ( 0.2303778133088964, 0.7148465705529154, 0 . 6308807679398587, -0.0279837694168599, -0.1870348117190931, 0.0308413818355607, 0.0328830116668852, - 0.0105974017850690 }; **/1 Filter from Eera Simoncelli ' s PhD thesis -- us ed in Edward Adelson's EPIC wavelet coder**  // These are probably the filter coefficients used in Shapiro's EZW paper Real AdelsonCoeffs[] = ( 0.028220367, -0.060394127, -0.07388188, 0.41394752, 0.7984298, 0.41394752, -0.07388188, -0 . 060394127, 0 . 028220367 ); *II 7/9 Filter from M. Antonini, M. Barlaud, P. Mathieu, and* **/1 I. Daubechies, -Image coding using wavelet transform-, IEEE**  // Transactions on Image Processing", Vol. pp. 205-220, 1992. Real AntoniniSynthesis  $|$  =  $(-6.453888262893856e-02,$ -4 .068941760955867e-02, 4.180922732222124e-01, 7.884856164055651e- 01, 4.180922732222124e-01, -4.068941760955867e-02,  $-6.453888262893856e-02$  }; Real AntoniniAnalysis[] = ( 3.782845550699535e-02, - 2 . 384946501937986e-02, - 1.106244044184226e- 01, 3.774028556126536e- 01, 8.526986790094022e- 01, 3.774028556126537e-01, -1.106244044184226e-01, - 2.38494650193 7986e-02, 3.782845550699535e-02 }; // Unpublished 18/10 filter from Villasenor's group Real Villa1810Synthesis  $(j) = (9.544158682436510e-04,$ -2.727196296995984e-06 , -9.452462998353147e-03, - 2 . 528037293949898e- 03, 3.083373438534281e-02 , ); -1.37 6513483818621e-02, -8.566118833165798e-02,

1.533685405569902e- 01, 6 . 233596410344172e-01, 6 . 233596410344158e- 01, 1 . 633685405569888e - 01, - 8 . 566118833165885e-02, -1 .376513483818652e-02, 3.08337343853426 7e -02, -2.528037293949898e-03 , - 9 . 452462998353147e - 03, - 2 .72 719629699598 4e - 06,  $9.544158682436510e-04$  ; Real Villa1810Analysis  $($  | =  $($  2.885256501123136e-02, 8.244478227504624e-05,  $-1.575264469076351e-01$ , 7.679048884591438e- 02 , 7.589077294537618e-01, 7 . 58907729453761ge- 01, 7 . 679048884691436e - 02, -1.575264469076351e-01, 8.244478227504624e-05, 2 . 885256501123136e-02) ; **1/ Filters from Chris Brislawn's tutorial code**   $Real BritishWnAnalysis$   $( | = | 0.037828455506995, -0.023849465019380,$ -0.110624404418423, 0 . 377402855612654, 0 . 852698679009403, 0 . 377402855612654, -0 . 110624404418423, -0.023849465019380, 0 . 037828455506995) ; Real BrislawnSynthesis  $[$   $] =$   $($  -0.064538882628938, -0.040689417609558, 0.418092273222212, 0.788485616405664, 0 .418092273222212, - 0.040689417609558, -0 . 064538882628938}; Real Brislawn2Analysis  $j$  =  $( 0.026913419, -0.032303352,$ -0.241109818, 0.054100420, 0.899506092, 0.899506092, 0 . 054100420, - 0 . 241109818, - 0.032303352, 0.026913419}; Real Brislawn2Synthesis  $[$   $] =$   $($  0.019843545, 0.023817599, -0 . 023257840, 0.145570740, 0.541132748, 0.541132748, 0.145570740, -0.023257840, 0.023817599, 0.019843545 }; **II Filters from J. Villasenor , B. Belzer, <sup>J</sup> . Liao, ·Wavelet Filter II Evaluation for Image Compression.- IEEE Transactions on Image**  */1* Processing, Vol. 2, pp. 1053 - 1060, August 1995. Real Villa1Analysis  $| \cdot |$  = ( 3.782845550699535e- 02, -2.384946501937986e-02, -1.106244044184226e-01, 3.774028556126536e- 01, 8.526986790094022e- 01, 3.774028556126537e-Ol, -1.10624404 4184226e-Ol, -2.384946501937986e-02, 3 . 782845550699535e-02 Real Villa1Synthesis  $|$  =  $\langle$ 



## **filter.cpp**<br>
Real \*anLow, int anLowSize, int anLowFirst,

```
FilterSet Haar (FALSE, HaarCoeffs, 2, 0);<br>FilterSet Daub4 (FALSE, Daub4Coeffs, 4, 0);
FilterSet Daub4 (FALSE, Daub4Coeffs, 4, 0);<br>FilterSet Daub6 (FALSE, Daub6Coeffs, 6, 0);
FilterSet Daub6 (FALSE, Daub6Coeffs, 6, 0);<br>FilterSet Daub8 (FALSE, Daub8Coeffs, 8, 0);
                   FALSE, Daub8Coeffs, 8, 0);<br>(TRUE, AntoniniAnalysis, 9, -4,
Filtext Antonini (TRUE, AntoniniAnalysis,
                           AntoniniSynthesis, 7, -3):
FilterSet Villa1810 (TRUE, Villa1810Analysis, 10, -4,
                           Villa1810Synthesis, 18, -8);<br>AdelsonCoeffs, 9, -4);
FilterSet Adelson (TRUE, AdelsonCoeffs, 9, -4)<br>PilterSet Brislawn (TRUE, BrislawnAnalysis, 9, -4,
FilterSet Brislawn (TRUE,
                           BrislawnSynthesis, 7, -3);
FilterSet Brislawn2 (TRUE, Brislawn2Analysis, 10, -4,
                           Brislawn2Synthesis, 10, -4) ;
FilterSet Villa1 (TRUE, Villa1Analysis, 9, -4, Villa1Synthesis, 7, -3);
FilterSet Villa2 (TRUE, Villa2Analysis, 13, -6, Villa2Synthesis, 11, -5);
FilterSet Villa3 (TRUE, Villa3Analysis, 6, -2, Villa3Synthesis, 10, -4);
FilterSet Villa4 (TRUE, Villa4Analysis, 5, -2, Villa4Synthesis, 3, -1);
FilterSet VillaS (TRUE, VillaSAnalysis, 2, 0, VillaSSynthesis, 6, -2 ) ; 
FilterSet Villa6 (TRUE, Villa6Analysis, 9, -4, Villa6Synthesis, 3, -1);
Filte rSet Odegard (TRUE, OdegardAnalysis, 9, -4, OdegardSynthesis, 7, - 3) ; 
I I Destructor 
Filter, ,-Filter () 
 { 
 if (coeff <math>\rightarrow</math> NULL)delete [J coeff; 
/void Filter::init (int filterSize, int filterFirst, Real *data)
 { 
   si ze = filterSize; 
   firstlndex = filterFirst; 
   center = - firstlndex; 
   coeff = new Real [size]; 
   if (data != NULL) (for (int I = 0; i < size; i+1coeff[i] = dista[i];else { 
     for (int i = 0; i \leq 3ize; i \neq i)
       coeff[1] = 0;void Filter::copy (const Filter& filter)
 { 
  if (coeff <math>! = NULL)delete [] coeff;
  init (filter.size, filter.firstIndex, filter.coeff);
FilterSet: :FilterSet (inc symmetric ,
```

```
Real *synLow, int synLowSize, int synLowFirst) :
           symmetric (symmetric) 
inc i, sign; 
analysisLow new Filter (anLowSize , anLowFirst, an Low) ; 
II If no synthesis c oeffs are given, assume wavelet is orthogonal 
if (synLow == NULL)
  synthesisLow = new Filter ('analysisLow):II For orthogonal wavelets, compute the high pass filter using
  // the relation q_n = (-1)^n n h<sub>1</sub>(1-n)^{n*}II (or equivalently q_1(1-n) = (-1)^{n}(1-n) h_n<sup>n*</sup>)
  anal ysisHigh new Filter (analysisLow->size, 2 - analysisLow->size 
                              analysisLow->firstlndex); 
  II Compute (-l)'(l-n) for first n 
  if (analysisLow->firstlndex % 2) 
   sign = 1;
  else sign = - 1 i
  for (i = 0; i < analysisLow->size; i++) (analysisHigh- >coeff[l - i - analysisLow->firstlndex -
                         analysisHigh->firstlndex] 
      sign'" analysisLow->coeff[iJ; 
    assert (1 - i - analysisLow->firstlndex -
                         analysisHigh->firstlndey. >= 0); 
    assert (1 - i - analysisLow->firstlndex -
                         analysisHigh->firstlndex < analysisHigh->size); 
    sign *= -1; 
  II Copy the high pass analysis filter to the synthesis filter 
  synchesisHigh = new Filter (*analysisHigh); 
else ( 
  II If separate synthesis coeffs given, assume biorthogonal 
  svnthesisLow = new Filter (synLowSize, synLowFirst, synLow);
  II For 
orthogonal wavelets, compute the high frequency filter using 
  // the relation g_n = (-1)^n complement (h^-(1-n)) and
  II// (or equivalently g_{-}(1-n) = (-1)^{n}(1-n) complement (h_{-}n))
                 g^- n = (-1) 'n complement (h_(1-n))
  analysisHigh new Filter (synthesisLow- >size, 2 - synthesisLow->size -
                       synthesisLow->firstIndex) ; 
  II Compute (-l)'(l-n) for first n 
  if (synthesisLow->firstlndex % 2) 
   slan = 1 :
  else sign = -1;
  for (i = 0; i < synthesisLow->size; i++) \uparrowanalysisHigh->coeff[l - i - synthesisLow->firstIndex -
                        a nalysisHigh->fi rstlndexj 
      sign * synthesisLow->coeff[il; 
    assert (1 - i - synthesisLow- >f irstIndex -
                         analysisHigh->firstlndey. >= 0 ) ;
```
assert (1 - i - synthesisLow->firstlndex -

### filter.cpp

 $14 - 22 - 22 - 22 - 22 - 22 - 22$ 

**sign \* = -1 ;**  analysisHigh->firscIndey. < analysisHigh->size);

```
synthesisHigh = new Filter
                       (analysisLow- >size, 2 - analysisLow- >size -
                       analysisLow->firstIndex);
```

```
// Compute (-1)^{A}(1-n) for first n
if (analysisLow- >firstIndex % 2) 
  sign = 1;
else sign = -1;
```

```
for (i = 0; i < analysisLow->size; i++) {
 synchesisHigh- >coeff(l - i - analysisLow->fi r stlndex -
                      synthesisHigh->firstIndex] =
   sign * analysisLow->coeff[i];
 assert (1 - i - analysisLow->firstlndex -
                      synthesisHigh->firstIndex >= 0); 
 assert (1 - i- analysisLow->firstlndex -
                      synthesisHigh->firstlndex < synthesisHigh->size);
```

```
sign *= -1;
```
FilterSet::FilterSet (const FilterSet& filterset) (

```
copy (filtersec);
```
FilterSet: : "FilterSet () (

```
delete analysis Low; 
delete analysisHigh;
delete synthesisLow; 
delete synchesisHigh;
```
**FilcerSet& FilcerSet: :operator= (const FilterSet filterset)** 

```
delete analysisLow; 
delete analysisHigh; 
delete synthesisLow; 
delete synthesisHigh; 
copy (filterset); 
return *this;
```
 $\overline{\mathcal{L}}$ 

(

void FilterSet::copy (const FilterSet& filterset)

**symmetric = filcerset.symmetric; analysisLow = new Filter** (~(filterset.analysisLow)); analysisHigh = new Filter  $(*(filterest.analysisHigh))$  ; **synthesisLow; new Filter (\*{filterset.synthesisLow));**   $synthesisHigh = new Filter (*(filterset.synthesisHigh));$ 



### global.cpp

(

 $\Delta$ 

~include <stdarg.h> **.include <stdio . h>**  linclude <stdlib.h> **.include <new .h>**  linclude <math.h> **#inc lude <assert.h>**  linclude "global.hh" lifdef DEBUG static FILE \*debug\_file; **static inc debug\_file\_open**  FALSE; **iendi f** 

### II function called when out of memory put a debugger breakpoint here

```
II if trying to locate c ause of Out of memory error
```

```
void no_more_memory ()
```

```
error ("Out of memory");
```
**1/ Initialize system-level things** 

#### void init()

(

```
1/ Call no_more_memory when unable to malloc
set_new_handler (no_more_memory);
```
#ifdef DEBUG debug\_file = fopen  $('delay.log', 'w+'');$  $debug_file\_open = (debug_file != NULL);$ **#endif** 

**II Close down system- level stuff** 

void shut\_down ()

```
( 
hfdef DEBUG 
   fclose (debug_file);
   debug_fi le_open = FALSE; 
lendif 
)
```
**volatile void error (char \*format,** . . . )

**va\_list list;** 

**va\_start (list, format);** 

printf ("Error: "); vprintf (format, list); **va\_end (list);**  <sup>p</sup> rinc f ( " \n " );

#ifdef DEBUG if (debug\_file\_open)

```
fprintf (debug_ file, "Error, " ); 
   vfprintf (debug_ file, format, list); 
   fprintf (debug file, "\n");
    fflush (debug_file); 
lendif 
  assert(O) ; 
void warning {char *forrnat, . .. } 
  va_ list l ist; 
  va_start (list, format) ; 
#ifdef DEBUG
  if (debug_file_open) (
   fprintf (debug_file, "Warning: ");
   vfprintf (debug_ file, format, list); 
   fprintf (debug_file, '\n');
   fflush (debug_file) ;
lendif 
  printf ( "Warning: ");
  vprintf (format, list) i
  va_end (list); 
  printf ('\n');
```
### pgm2raw.cpp

```
'include <stdio.h> 
#include <stdlib.h>
#include <string.h>
#include <math.h>
l include <iostream.h> 
l include "global . hh" 
l include "image.hh " 
int main (inc argc , char **argv) 
  Image *image; 
  char * program = argv[0];if (argc != 3) 
fprintf (stderr, 
           "Convert an image in pbm/pgm format to raw pixel format\n " ); 
    fprintf (stderr, 
           · Usage : %s (pgm image name) [raw image name]\n - , program); 
    return 1-
  \Deltaimage = new Image (argv(1)) ;
  image - >saveRaw (argv[2J); 
  return 0;
```
quantize.cpp



//Ouantize.cpp  $1*$  . The contribution of the contribution of the contribution of the contribution of  $\mathcal{O}(n^2)$ #include <stdig.h> #include <stdlib.h> #include <math.h> Winclude <iostream.h> Winclude <iomanip.h> \*include "global.hh" #include "quant.hh" Ouantizer::Ouantizer (ErrorMetric "err) : err (err)  $data = NULL$  $nData = 0 +$  $max = min = sum = sumSG = mean = var = 0$ :  $initialDist = 0$ : void Ouantizer: : getStats ()  $max = -Max$   $= 1$ . min = MaxReal:  $sum = sumSa = 0$ : for  $(int i = 0; 1 < nData; 1++)$  $if (data[i] < min)$  $min = data[i];$ if (data[i] > max)  $max = data[i]$  $sum \rightarrow= data[i];$  $sumSq += square(data[i])$ ; mean =  $\sin$  / (Real)nData:  $var = sumSq / (Real)nbata - square (mean)$ : UniformOuant::UniformOuant (MonoLaverCoder \*entropy, int paramPrecision, int zeroCenter, ErrorMetric "err) : Quantizer (err), entropy (entropy), paramPrecision (paramPrecision), zeroCenter (zeroCenter) void UniformOuant::setDataEncode (Real "newData, int newNData)  $data = newData$ : nData = newNData; getStats (); if (zeroCenter) (  $max = fabs(max) > fabs(min)$  ?  $fabs(max)$  :  $fabs(min)$ :  $min = -max$ imax = realToInt (max, paramPrecision);  $cmax = intToReal (inax, paramPrecision):$ // Make sure cmax  $>$  max and -cmax <= min

while  $(\text{max} \leq \text{max})$ 

```
qmax = intToReal (++imax, paramPrecision);
 if (zeroCenter) (
   imin = -imax:
   min = -max;
 i else iimin = realToInt (min, paramPrecision);
   cmin = intToReal (imin, paramPrecision);
   // Make sure \dim <= \minwhile (\text{cmin} > \text{min})qmin = intToReal (--imin, paramPrecision);
 imean = realToInt (mean, paramPrecision);
 omean = intToReal (imean, paramPrecision);
 initialDist = 0:
 if (zeroCenter)
   for \{int\} = 0; \{x \in \mathbb{R}^n : x \neq 0\}initialDist += ("err) (datalil);
 a1cafor \{int i = 0 : i < nData: i++)\}initialDist += (*err) (data[i] - qmean) :
void UniformOuant: : setDataDecode (Real *newData, int newNData,
                              int imax, int imin, int imean)
 data = newData:
 nData = newNData;if (imin < imax) (
   cmax = intToReal (imin, paramPrecision);
   if (zeroCenter) (
     min = -maxi else icmin = intToReal (imin, paramPrecision):
   gmean = intToReal (imean, paramPrecision);
void UniformQuant::getRateDist (int precision, Real minStepSize,
                             Real &rate, Real &dist)
 if (precision > 01 (
   const int nSteps = (1<precision)-1;
   const Real stepSize = (\text{max-min}) / (\text{Real})nSteps;
   const Real recipStepSize = 1.0/stepSize:
   if (stepSize < minStepSize) (
     rate = MaxReal:dist = MaxReal:return:
   entropy->setNSym (nSteps);
   rate = dist = 0:
   for (int i = 0; i < nData; i++) (
```
#### quantize.cpp

```
int symbol = (int) ((data[i] -qmin) *recipStepsize);
      assert (symbol < nSteps && symbol >= OJ ; 
      rate += entropy- >cost (symbol, TRUE); 
      Real reconstruct = qmin + ((Real)symbol + 0.5) * stepSize;
      dist += ("err) (data[i]-reconstruct);
 else ( 
    rate = dist = 0 ; 
    if (zeroCenter) ( 
     for (int i = 0; i < nData; i + 1 (dist += (*err) (data[i]];
   else ( 
      for (int i = 0; i < nData; i++) { 
        dist \div = (*err) (data[i] - amean):voi d UniformQuant: :quantize (Encoder *encoder, int precision) 
( 
 if (precision> 0) ( 
   const int nSteps = (1<corecision)-1;
    const Real stepSize = (\text{dmax-cmin}) /(Real)nSteps;
    canst Real recipStepSize = 1.0/scepSize; 
    entropy->setNSym (nSteps); 
    for (int i = 0; i < nData; i \rightarrow 0int symbol = (int) ((data[i] -qmin)*recipsLength);\text{assert} (symbol < \text{nSteps} && symbol >= 0) ;
      entropy->write (encoder, symbol, TRUE); 
voi d UniformQuant::dequantize (Decoder *decoder, int precision) 
( 
  if (precision > 0) (
   canst int nSteps = (l«precision) - l; 
    const Real stepSize = (\text{max-cmin})/(\text{Real}) nSteps;
    int symbol; 
    entropy->setNSym (nSteps);
    for (int i = 0; i < nData; i++) (
      symbol = entropy->read (decoder, TRUE); 
      assert (symbol < nSteps 64 symbol >= 0);
      data[i] = qmin + ((Real)symbol + 0.5) * stepSize;else ( 
    for (Int i = 0; i < nData; i++) (
     data[i] = qmean;/ We are not a produced denote a low a state of the state and denote below.
void UniformQuant. ::writeHeader (Encoder *encoder, int precision)
```

```
encoder->writeNonneg (precision);
 if (precision > 0) (
   encoder- >writelnt (imax); 
   if (! zeroCenter) 
     encoder - >writel nt (i min); 
 \lvert else lif (!zeroCenter) 
     encoder ->wri telnt (imean) ; 
void UniformQuant' ,readHeader (Decoder *decoder, int &precision) 
( 
 precision = decoder - \geq readNonneq () ;if (precision > 0) (
   imax = decoder->readInt ();
   qrnax = intToReal (imax , paramPrecision) i
   if (zeroCenter) 
     gmin = \text{-} \text{cmax:}else 
     imin decoder->readlnt (); 
     cmin = intToReal (imin, paramPrecision) ;
   \text{mean} = 0;
 else ( 
   if (! zeroCenter) 
     imean = decoder->readlnt () ; . 
     qrnean = intToReal (imean , paramPrecision) ; 
   else ( 
     \text{cmean} = 0;
     i mean = i realToInt (qmean, paramPrecision) ;
   qmax = qmin = qmean;
void UniformQuant : :setParams (int newParamPrecision, Real newMax, 
                           Real newMin, Real newMean) 
  paramPrecision 
newParamPrecision; 
 qmax = newMax; 
 qrnin = newMin; 
 qrnean = newMean; 
LayerQuant"LayerQuant (MultiLayerCoder 'entropy, int paramPrecision, 
                       int signedSym, int nLayers, ErrorMetric ~err) 
 Quantizer (err), entropy (entropy), paramPrecision (paramPrecision),
 signedSym (signedSym), nLayers (nLayers)
 currentLayer = -1; 
 layerRate = new Real InLayers); 
 layerDist = new Real [nlayers];context = NULL; 
 residual = NULL;
```

```
quantize.cpp
```
void LayerQuant::setDataDecode (Real \*newData, int newNData, int imax, **int imin , int imean)** 

```
LaverOuant: : "LaverOuant ()
( 
 delete I) layerRate; 
 delete [] layerDist;
 if (context != NULL) 
   delete II context;
 if lresidual != NULL)
   delete [] residual;
void LayerQuant : :setDataEncode (Real *newData, i nc newNData) 
( 
 data:; newData: 
 nData :; newNData: 
 getStats ();
 if (signedSym) 
   max = fabs(max) > fabs(min) ? fabs(max) : fabs(min);
   min = -max:
 imax = realTolnt (max, paramPrecision); 
 qrnax = intToReal (imax, paramPrecision); 
 / / Make sure qrnax >= max and - qmay. <= min 
 while (qrnax < max) 
   qrnax = intToReal (++imax, paramPrecision); 
 if (signedSym) ( 
   imin = -imax;
   qmin = -qmax;
 else 
   imin realTolnt (min , paramPrecision); 
   qmin intToReal (imin, paramPrecision); 
   II Hake sure qmin <:::: min 
   while (qmin > min)qrnin = intToReal (--imin , paramPrecision); 
 if (signedSym) 
   threshold = qmax/2.0; 
 else 
   threshold = (\text{max - qmin})/2.0;currentLayer = -1;
 if (context i = NULL)delete () context;
 if (residual != NULL) 
   delete (1 residual;
 context;; new int [nDatal; 
 residual = new Real InData1;resetLayer (); 
 ini tialDist = 0; 
 for (int i = 0; i < nData; i++)initialDist += ("err)(residual[i]);entropy->reset ();
```

```
data = newData; 
 nData = newNData; 
 if (\text{imin} < \text{imax})qmax = intToReal (imin, paramPrecision);
   gmean = intToReal (irnean, pararnPrecision) ; 
   if (signedSym) ( 
     qmin = -qmax;threshold = \text{cmax}/2.0;
   else ( 
     qmin = intToReal (imin, pararnPrecision); 
     threshold = (\text{cmax} - \text{cmin})/2.0;
 currentLayer = -1 ; 
 if (context != NULL)
   delete [J context; 
 context = new i nt [nData] ; 
 if (signedSym) (
   for (int i = 0; i < nData; i++) { 
     data[i] = 0;context[i] = 0;
  else ( 
   for (int i = 0; i < nData; i \leftrightarrow i) {
     data[i] = qmean;context[i] = 0; 
 resetLayer (); 
  entropy->reset (); 
void LayerQuant: :getRateDist (int precision, Real minStepSize, 
                              Real &rate, Real &dist) 
 assert (precision <= nLayers) ; 
  Real currentRate 
0 ; 
  Real currentDist 
ini tialDist; 
 for (int i = 0; i < precision; i++) (
   II rates & distortions have been computed for layers up to currentLayer 
   if (i \leq currentLayer) (i \leq j \leq n)currentRate += layerRate[i];
     currentDist += layerDist[i}; 
   else ( 
     if (threshold > minStepSize) (
       quantizeLayer (NULL) ; 
        currentRate += laverRate[i];
        currentDist += layerDistlil; 
     else ( 
        layerRate[i) = MaxReal; 
        layerDist{iJ = - MaxReal; 
        currentRate = MaxReal;
```
Real deltaRate =  $0$ , deltaDist =  $0$ ;

```
currentDist 
MaxReal; 
rate currentRate;
```

```
dist = currentDist;
```
**void LayerQuant::quantize (Encoder \*encoder, int precision)** 

```
resetLayer I); 
entropy->reset () ;
```
I

(

```
for (\text{int } i = 0; i < \text{precision}; i++) (
  quantizeLayer (encoder);
```
**void LayerQuant. ::resetLayer ()** 

```
if (residual != NULL) 
  if (signedSym) ( 
   for (int i = 0 ;i < nData; i++) ( 
      residual[i] = data[i];context[i] = 0;
  else ( 
    li on the first layer we remove the mean 
    for (int i = 0; i < nData; i++) (
     residual[i] = data[i] - 0.5*(qmax+qmin);context[i] = 0;else 
  if (signedSym) 
    for lint i = 
:i < nData; i++) { 
      context[i] = 0;e lse { 
    lI on the first layer we remove the mean 
    for (int i = 0; i < nData; i++) (
     context[i] = 0;currentLayer = -1;
```

```
if (signedSym) 
 threshold = qrnax/2 . 0; 
else 
  threshold = \frac{1}{2}.0;
```
*II* deltaRate **bits required to code current. layer**  *II* deltaDist **reduction in distort ion from current layer** 

**void LayerQuant: :quantizeLayer (Encoder \*encoder)** 

```
const Real halfThreshold = 0.5 * threshold;
const Real threeHalvesThreshold = 1.5 * threshold;
```

```
int symbol; 
currentLayer++; 
II printf ('current layer = %d, threshold %g\n', currentLayer, threshold); 
if (signedSym) I 
  for (int i = 0; i < nData; i \leftrightarrow i)
    deltaDist -= (*err) (residual(i)); // subtract off old error<br>// Real oldResid = residual(i);
         Real oldResid = residual[i];
    if (context[i] == 0) (
      if (residual[i] > threshold)swmbol = 1:
         residual(i) -= threeHalvesThreshold; 
      \big) else if (residual(i) < -threshold) (
         symbol{1 = -1};
         residual Ii) += threeHalvesThreshold; 
       ) else (
         symbol{1 = 0};
     else I 
      if (residual(i) > 0) (
         symbol1 = 1;
         residual Ii) -= half Threshold; 
      else I 
         swmbol = 0;
         residual [i) += half Threshold ; 
    deltaDist *= (*err) (residual(i)) ; // add in new error
    deltaRate entropy->write (encoder, symbol, TRUE , currentLayer, 
                                   context[i]);
    content[i] = 2*context[i] + symbol;else 
  for (int i = 0; i < nData: i++) ( 
    deltaDist -= (*err) (residual | i]); // subtract off old error
    if (residual[i] > 0) (
      symbol1 = 1;
      residual [i) -= half Threshold; 
     else I 
      swmbol = 0;
      residual Ii) += half Threshold; 
    deltaDist (*err) (residualli)); II add in new error 
    delt.aRate += entropy->write (encoder, symbol, TRUE, currentLayer, 
                                   context[i]);
    contextli) = 2*context[i) + symbol; 
threshold *= 0.5;layerRate(currentLayer) = deltaRate;
layerDist[currentLayer] = deltablist;
```
# quantize.cpp

void LaverOuant::dequantize (Decoder "decoder, int precision)

```
resetLaver ();
```

```
for \{int i = 0: i < precision; i++) (
 dequantizeLaver (decoder);
Y
```
void LaverOuant::dequantizeLaver (Decoder \*decoder)

int symbol;

currentLayer++;

```
if (signedSym) (
 for (\text{int } i = 0; i < n\text{Data}; i++) (
    symbol = entropy->read (decoder, TRUE, currentLayer,
                               context[i]);
```

```
11int oldData = data[i];
if (context[i] == 0)
  data[i] += 1.5*threshold * symbol;
else
  data[i] \div (symbol - 0.5) * threshold;
```

```
context[i] = 2*context[i] + symbol;) else (
 for (int i = 0: i < nData: i++) (
   symbol = entropy->read (decoder, TRUE, currentLaver, context(i));
```

```
data[i] += (symbol - 0.5) * threshold;
context[i] = 2*context[i] + symbol:
```

```
threshold *= 0.5;
```
void LayerQuant::writeHeader (Encoder \*encoder, int precision)

```
encoder->writeNonneg (precision);
```

```
if (precision > 0) {
 encoder->writeInt (imax);
 if (!signedSym)
   encoder->writeInt (imin);
) else (if (!signedSym)
   encoder->writeInt (imean):
```
precision = decoder->readNonneg ();

```
void LayerQuant:: readHeader (Decoder *decoder, int &precision)
```

```
if (precision > 0) (
   imax = decoder->readInt ();
   gmax = intToReal (imax, paramPrecision);
   if (signedSym) (
    cmin = -cmax;\text{cmean} = 0;
   j else (i)imin = decoder->readInt ();
    cmin = intToReal (imin, paramPrecision);
    \text{mean} = 0.5 * (\text{max} * \text{min}):1 else (if (isignedSym) (
    imes n = decoder->readInt ();
     cmean = intToReal (imean, paramPrecision);
  I else f
     \text{mean} = 0;
     imean = realToInt (qmean, paramPrecision);
   qmax = qmin = qmean;
```
void LayerQuant::setParams (int newParamPrecision, Real newMax, Real newMin, Real newMean)

```
paramPrecision = newParamPrecision:
\text{max} = newMax:
min = newMin:
dmean = newMean:
```




#include <stdio.h> #include <stdlib.h> #include <string.h> #include <math.h> #include <iostream.h> #include "global.hh" #include "image.hh" int main (int arge, char \*\*argv)  $\left($ Image \*image: char \*program =  $argv[0]$ ; if  $\text{trace} := 5)$  { fprintf (stderr, "Convert an image in pbm/pgm format to raw pixel format\n"): fprintf (stderr, 'Usage: %s [raw image name] [height] [width] [pgm image name] \n', program); return 1;  $\mathcal{F}$ int hsize =  $atoi(argv[3])$ ; int vsize =  $atoi(arav[2])$ ; image = new Image (argv[1], hsize, vsize); image->savePGM (argv[4]); return 0; 

# wavelet.cpp

#include <stdio.h> Winclude <stdarg.h> #include <stdlib.h> tinclude <math.h> #include cassert.h> \*include "global.hh" #include 'image.hh' \*include \*wavelet.hh\* Wavelet::Wavelet (FilterSet \*filterset) analysisLow = filterset->analysisLow: analysisHigh = filterset->analysisHigh: synthesisLow = filterset->synthesisLow; synthesisHigh = filterset->synthesisHigh; symmetric = filterset->symmetric; // amount of space to leave for padding vectors for symmetric extensions npad = max(analysisLow->size, analysisHigh->size); Wavelet: : "Wavelet () void Wavelet::transform1d (Real \*input, Real \*output, int size. int nsteps, int sym ext) int i: int currentIndex =  $0:$ Real \*data[2]; int lowSize = size, highSize: if  $l$ svm ext ==  $-1$ )  $swm$  ext =  $swmmetric$ : data [0] = new Real [2\*npad+size]; data  $111$  = new Real  $[2*$ npad+sizel; for  $(i = 0; i <$  size;  $i++)$  $data|currentIndex|$   $(npad+1) = input[i];$ while  $(nsteps--)$  ( if (lowSize <= 2 && symmetric == 1) { warning ("Reduce # of transform steps or increase signal size") ; warning (\* or switch to periodic extension\*); error ('Low pass subband is too small'): // Transform printf ("transforming, size = %d\n", lowSize); transform\_step (data[currentIndex], data[1-currentIndex],

lowSize, sym\_ext);

```
highSize = lowSize/2;
    lowSize = (lowSize + 1)/2:
    // Copy high-pass data to output signal
    nony (data (1-currentIndex) = noad + lowSize, output +
          lowSize, highSize):
    for (i = 0: i < 1owSize+highSize: i++)
      printf ("$5.2f", data[1-currentIndex][npad+i]);
    printf ("\n\n");
    // Now pass low-pass data (first 1/2 of signal) back to
    // transform routine
    currentIndex = 1 - currentIndex:
  // Copy low-pass data to output signal
  copy (data[currentIndex] + npad, output, lowSize);
  delete [] data [1]:
  delete [] data [0]:
void Wavelet:: invert1d (Real *input, Real *output, int size,
    int nsteps, int sym ext)
  int is
  int current<br>Index = 0:
  Real *data(2):
  if (svm ext = -1)swm ext = symmetric;
  int *lowSize = new int (nsteps):
  int *highSize = new int (nsteps):
  10WSize[0] = (size+1)/2:
  highSize[0] = size/2;
  for (i = 1; i < nsteps; i++) (
    lowSize[i] = (lowSize[i-1] + 1)/2:
    hichSize[i] = lowSize[i-1]/2;
  data [0] = new Real [2*npad+size);
  data |1| = new Real [2*npad+size]:
  copy (input, data(currentIndex)+npad, lowSize(nsteps-1));
  while (nsteps--) (
    // grab the next high-pass component
    copy (input + lowSize[nsteps],
          datalcurrentIndex]+npad+lowSize[nsteps], highSize[nsteps]);
    // Combine low-pass data (first 1/2*n of signal) with high-pass
    // data (next 1/2^n of signal) to get higher resolution low-pass data
    invert_step (data(currentIndex), data[1-currentIndex],
                 lowSize[nsteps]+highSize[nsteps], sym_ext];
```
#### wavelet.cpp

*II* Now pass low-pass data (first *1/2* of signal) back to **II transform routine currentlndex ;;; 1 - currentlndex;** 

*II* Copy inverted signal to output signal copy (datalcurrentlndexl+npad, output, size);

delete II highSize; delete I) lowSize;

delete () data (1); delete [] data [0];

**void Wavelet : :cransform2d (Real \*input, Real \*output, inc hsize, inc vsize, int nsteps, inc sym\_ext)** 

int j; **inc hLowSize**  hsize, hHighSize; **inc vLowSize vsi ze, vHighSize;** 

if  $(sym\_ext == -1)$  $sym$  ext = *symmetric*;

**Real Ycemp\_in = new Real [2\*npad+max(hsize,vsize)];**  Real \*temp\_out = new Real  $[2*npad+max(hsize,wise)]$ ;

**copy (input, output, hsize\*vsize);** 

while (nsteps--) ( if ((hLowSize  $\leq$  2 || vLowSize  $\leq$  2) && sym\_ext == 1) { **warning (-Reduce # of transform steps or increase signa <sup>l</sup>***size- );*  **warning (- or switch to periodic extension - );**  error ("Low pass subband is too small");

 $for (i = 0; j < v$ LowSize;  $j++)$  { *II* Copy row j to data array copy (output+(j\*hsize) , temp\_in+npad, hLowSize);

transform\_step (temp\_in, temp\_out, hLowSize, sym\_ext);

*II* Copy back to image copy (temp\_out+npad, output+(j\*hsize), hLowSize);

for  $(j = 0; j < h$ LowSize;  $j \rightarrow ()$  ( *II* Copy column j to data array **copy (OUCput+j, hsize, cernp\_in+npad, vLowSize);** 

**II Convolve with low and high pass filters transform-step (temp\_in, temp\_out, vLowSize, sym\_ext);** 

**copy** (temp out~npad, output~j, **hsize, vLowSize ); II Now convolve low-pass portion again**  hHighSize = hLowSize/2; hLowSize = (hLowSize+1) *12;*  vHighSize = vLowSize/2; vLowSize = (vLowSize+1) *12;*  delete II temp\_out; delete (1 temp in: **void Wavelet::invert2d (Real 'input, Real 'output, int hsize, int vsize, int nsteps, int sym\_ext) int i, j;**  if  $(svm.ext = -1)$ **sym\_ext = symmetric; int \*hLowSize = new int Insteps], \*hHighSize = new int Insteps); int \*vLowSize = new int Insteps], \*vHighSize = new int Insteps];**  hLowSizelOI = *(hsize+1)/2;*  hHighSizelOI = *hsize/2;*  vLowSizelOI = *(vsize+ 1) / 2;*  vHighSizelOI = *vsize/2;*  **for (i = 1; i < nsteps; i++) (**   $h$ LowSize[i] =  $(h$ LowSize[i-11+1)/2;  $hHighSize[i] = hLowSize[i-1]/2$ ;  $v$ LowSize[i] =  $(v$ LowSize[i-1]+1]/2; vHighSizelil = vLowSizeli-11 *12;*  **Real \*temp\_in = new Real [2\*npad+rnay.(hsize,vsize)]; Real \*temp\_out = new Real (2\*npad+max(hsize,vsize)]; copy (input, output, hsize\*vsize);**  while (nsteps--) ( **II Do a reconstruction for each of the columns**  for  $(j = 0; j < h$ LowSize(nsteps]+hHighSize(nsteps];  $j++)$ *I* / Copy column j to data array copy (output+j, hsize, temp\_in+npad, vLowSizelnsteps]+vHighSizelnsteps]) ; *II* Combine low-pass data (first *<sup>1</sup> / 2'n* of signal) with high-pass *II* data (next *1/2'n* of signal) to get higher resolution low-pass data **invert\_ seep (temp\_ in, temp\_out,**  vLowS izeinsteps]+vHighSizelnsteps), **sym\_ex t) ;**  *II* Copy back to image c opy (temp\_out+npad, output+j, hsize,

vLowSizelnsteps] +vHighSizelnstepsl );

*II* Copy back to image

wavelet.cpp

```
1/ Now do a reconstruction pass for each row 
      for (j = 0; j < vLowSize[nsteps]+vHighSize[nsteps]; j++) (
         II Copy row j to data array 
         copy (output + (j*hsize), temp_in+npad,
               hLowSizelnstepsl+hHighSizelnstepsll; 
         II Combine low-pass data (first 1/2'n of signall with high-pass 
         II data (next 1/2' n of signal) to get higher resolution low-pass data 
         invert_ step (temp_in, temp_ out, 
                     hLowSize[nsteps]+hHighSize[nsteps], sym_ext);
         II Copy back to image 
         copy (temp_out+npad, output + (j*hsize),
               hLowSize(nstepsl+hHighSizelnstepsl) ; 
  delete (I hLoWSize; 
   delete II hHighSize; 
  delete II vLowSize; 
  delete () vHighSize;
   delete II temp_in: 
   delete II temp_out; 
void Wavelet: :transform_step (Real *input, Real *output, inc size, 
                              int syrn_ext) 
 inc i, j; 
 int lowSize = (size+1)/2; 
 inc left_ext, right_ext; 
 if (analysisLow->size %2) 
   II odd filter length 
   left\_ext = right\_ext = 1;else ( 
    left\_ext = right\_ext = 2;if (syrn_ext) 
   symmetric_extension (input, size, left_ext, right_ext, 1);
   periodic_extension (input, size); 
 II coarse detail 
  II xxxxxxxxxxxxxxxx --> HHHHHHHHGGGGGGGG 
  for (i = 0; i < 1owSize; i++) (
   output[quad+1] = 0.0;for (j = 0; j < analysisLow->size; j++)output Inpad+iJ +; 
        inputlnpad + 2*i + analysisLow->firstlndex + j] * 
        analysisLow->coeff{jl ; 
  for (i = lowSize; i < size; i++) 
                                                                                                \mathcal{L}
```
**else** 

 $output[npad+1] = 0.0;$ 

```
for (j = 0; j < analysisHigh->size; j++) (
     output (npad+iJ += 
       input[npad + 2*(i-lowSize) + analysisHigh-5firstIndex + j] *analysisHigh- >coeff(jl; 
void Wavelet::invert_ step (Real *input, Real *output, int size, int syrn_ ext) 
  int i, j; 
  int left_ext, right_ext, symmetry; 
  II amount of low and high pass -- if odd \ast of values, extra will be
  11 low pass
  int lowSize = (size+1)/2, highSize = size/2; 
   symmetry = 1;
   if (analysisLow->size \$ 2 == 0) (
    II even length filter -- do (2, X) extension 
    left_ext = 2;
   else ( 
     // odd length filter -- do (1, X) extension
     left\_ext = 1;
   if (size % 2 == 0) (
    II even length signal -- do (X, 2) extension 
     right ext = 2;
   else ( 
     1/ odd length signal -- do (X, 1) extension
     right_ext = 1; 
   Real *temp = new Real [2*npad+lowSiz ej; 
   for (i = 0; j < 1owSize; i+1 (
    \text{temp}[\text{npad}+i]=\text{input}[\text{npad}+i];if Isyrn_ext) 
    symmetric_extension (temp, lowSize, left_ext, right_ext, symmetry); 
   else 
    periodic_extension (temp, lowSize) ; 
   II coarse detail 
   // HHHHHHHGGGGGGGGGGG --> xxxxxxxxxxxxxxxxx
   for (i = 0; i < 2*npad+size; i++) 
    output[i] = 0.0;int firstIndex = synthesisLow->firstIndex; 
   int last.Index = synthesisLow->size - 1 + firstIndex; 
   for (i = -lastIndex/2; i \leq (size-1-firstIndex)/2; i++)for (j = 0; j < synthesisLow->size; j++) (
       output(npad + 2*i + firstIndex + j +=
         templnpad+il • synthesisLow->coeffljl; 
   left_ext 2; 
   if (analysisLow->size k 2 == 0) (
```

```
// even length filters
 right ext = (size 2 == 0) 7 2 : 1:
 symmetry = -1:
) else (// odd length filters
 right ext = (size % 2 == 0) ? 1 : 2;
 symmetry = 1:
```

```
for (i = 0; i < highSize; i \leftrightarrow j (
 temp(npad+1) = input(npad+lowSize+1);
```
if (sym ext) symmetric\_extension (temp, highSize, left\_ext, right\_ext, symmetry); else

```
periodic_extension (temp, highSize);
```

```
firstIndex = synthesisHigh->firstIndex;
lastIndex = synthesisHigh-ssize - 1 + firstIndex;
```

```
for (i = -lastIndex/2; i \leq (size-1-firstIndex)/2; i++) (
  for (j = 0; j < synthesisHigh->size; j++) (
    output[npad + 2*1 + firstIndex + j] +=
      temp[npad+i] * synthesisHigh->coeff[j];
```

```
delete () temp;
```
void Wavelet: symmetric\_extension (Real \*output, int size, int left ext, int right\_ext, int symmetry)

```
int if
int first = npad, last = npad + size-1:
```

```
if (symmetry == -1) (
 if (left ext = 1)
   output[-first] = 0;if (right\_ext == 1)output[++last] = 0;int originalFirst = first;
int originalLast = last;
int originalSize = originalLast-originalFirst+1;
int period = 2 * (last - first + 1) - (left_ext == 1) - (right_ext == 1);
```

```
if (left\_ext == 2)output[--first] = symmetry*output[originalFirst];
if (right\_ext == 2)output[++last] = symmetry*output[originalLast];
```
int nextend = min (originalSize-2, first); for  $(1 = 0; 1 <$  nextend;  $1++1$  ( output[--first] = symmetry\*output[originalFirst+1+i];

```
while [t]irst > 0) (
 first--:
 output[first] = output[first+period];
```
 $nextend = min$  (originalSize-2, 2\*npad+size-1 - last); for  $(i = 0; i <$  nextend;  $i++)$  ( output[++last] = symmetry\*output[originalLast-1-i];

while (last < 2\*npad+size-1) (  $last++$  $outputu[last] = output[last-period];$ 

void Wavelet::periodic\_extension (Real \*output, int size) int first = npad, last = npad + size-1;

while (first  $> 0$ ) ( first- $output[first] = output[first+size];$ 

while  $\{last < 2^* \text{npad} + size - 1\}$  (  $last++$ output(last) = output(last-size);

## transfor.cpp

```
\mathbf{H} , and \mathbf{H} is a set of the set o
                                                                                    #include <stdio.h>
#include <stdlib.h>
                                                                                    void WaveletTransform::init ()
#include <math.h>
#include 'trans.hh'
                                                                                      int_1value = new Real (hsize*vsize):
WaveletTransform: : WaveletTransform (Wavelet *wavelet, Image *image,
                                                                                      nSubbands = 3 * nsteps + 1:int nsteps, int symmetric) :
                                                                                      subbandSize = new int (nSubbands):wavelet (wavelet), nsteps(nsteps), symmetric(symmetric)
                                                                                      subbandE size = new int [nSubbands]:subbandVsize = new int [nSubbands]:subbandPtr = new Real* [nSubbands]:value = NULLif (image \mid = NULL) (
                                                                                      int *lowHsize = new int (nsteps) :
   hsize = image \rightarrow hsizeint *lowUsing = new int Instead:vsize = image->vsize:int *highHsize = new int [nsteps];
   transform (image, wavelet, nsteps, symmetric):
                                                                                      int *highVsize = new int (nsteps):
 i else ihsize = vsize = 0;
                                                                                      lowRsize-Insteps-11 = (hsize+1)/2:lowVsizelnsteps-11 = (vsize+1)/2:
                                                                                       highHsizeInsteps-11 = hsize/2:
                                                                                      highVsize[nsteps-1] = vsize/2;
for (i = nsteps-2; i > = 0; i--) (
WaveletTransform: WaveletTransform (Wavelet *wavelet, int hsize, int
                                                                                        lowHsize[i] = (lowHsize[i+1]+1)/2;
                               vsize, int nsteps, int symmetric):
                                                                                        lowVsize[i] = (lowVsize[i+1]+1)/2:
                      hsize(hsize), vsize(vsize), wavelet
                                                                                        highHsize[i] = lowHsize[i+1]/2;(wavelet), nsteps(nsteps), symmetric(symmetric)
                                                                                        highVsize[i] = lowVsize[i+1]/2;
 nsteps = 0;
 symmetric = -1;
                                                                                       subbandPtr[0] = value;init ():
                                                                                       subbandHsize[0] = lowHsize[0];
 for (int i = 0: i < hsize*vsize: i++)subbandVsize[0] = lowVsize[0];value[i] = 0:
                                                                                       subbandSize[0] = subbandHsize[0]*subbandVsize[0];for (i = 0; i < nsteps; i++) (
subbandHsize(3*1+11) = highHsize[11]subbandVsize[3*1+1] = lowVsize[1];WaveletTransform::WaveletTransform (const WaveletTransform &t)
                                                                                        subbandHsize[3*1+2] = lowHsize[i];subbandVsize[3*1+2] = highVsize[1];wavelet = t.wavelet;subbandHsize[3*1+3] = highHsize[1];hsize = t.hsize:
                                                                                        subbandVsize(3*1+3) = highVsize(i);vsize = t.ysizensteps = t.nsteps;symmetric = t.symmetric;
                                                                                       for (i = 1; i < nSubbands; i++) (
                                                                                        subbandSize[i] = subbandHsize[i]*subbandVsize[i];
 if (t \cdot value == NULL) (
                                                                                        subbandPtr[i] = subbandPtr[i-1] + subbandSize[i-1];value = NULL\frac{1}{2} else (init():
                                                                                       delete (1 lowHsize:
   for (int i = 0; i < hsize*vsize: i++)delete [] lowVsize:
     value[i] = t.value[i];delete () highHsize:
                                                                                       delete || highVsize;
WaveletTransform:: "WaveletTransform ()
                                                                                    void WaveletTransform::freeAll ()
  freeAll ():
                                                                                     if (value != NULL) (
                                                                                       delete [] value;
                                                                                       delete [] subbandSize;
```
## transfor.cpp

delete () subbandHsize; delete () subbandVsize: delete [] subbandPtr:

 $freeAll()$ :

 $int()$ :

 $hsize = image \rightarrow hsize:$ 

 $vsize = image->vsize:$ 

wavelet = newWavelet:  $nsteps = steps;$ 

// linearize data mallatToLinear (temp); delete () temp;

 $symmetric = isSymmetric$ :

Real \*temp = new Real [hsize\*vsize];

Real \*temp = new Real (hsize\*vsize);

// put data in Mallat format

int \*lowHsize = new int (nsteps):

int \*lowVsize = new int (nsteps):

 $lowHsize(nsteps-1) = (hsize+1)/2$ ;

 $lowVsize(nsteps-1) = (vsize+1)/2;$ 

for  $(i =$  nsteps-2;  $i \gg 0$ ;  $i-j$  (

 $lowHsize[i] = (lowHsize[i+1]+1)/2$ ;

 $lowVsize[1] = (lowVsize[i+1]+1)/2$ :

linearToMallat (temp);

delete il temp:

 $int 1, 1, k$ ;

```
for (i = 0, i < subbandHsize[0]; i \leftrightarrow jsubbandPtr[0][j*subbandHsize[0]+i] =
       mallat(i*hsize+i);
 for (k = 0; k < nsteps: k++) (
   for (i = 0; j < subbandVsize[k*3+1]; (i++)for (i = 0; i < subbandHsize[k*3+1]; 1++)subbandPtr[k*3+1][j*subbandHsize[k*3+1]+1] =
         mallat[j"hsize+(lowHsize[k]+i)];
   for (i = 0; j < subbandVsize[k*3+2]; (i+1)for (i = 0; i < subbandHsize[k*3+21; i++)subbandPetr(k=3+2)(i*subbandHsize(k*3+21+i) =mallat[(lowVsize[k]+j)*hsize+i];
   for (j = 0; j < subbandVsize[k*3+3]; j++)
     for (i = 0; i < subbandHsize[k*3+3]; i++]
       subbandPtr[k*3+3|[i*subbandHsize[k*3+3|+i] =
         mallat((lowVsize(k)+j)*hsize+(lowHsize(k)+i));
 delete () lowHsize:
 delete [] lowVsize;
void WaveletTransform::linearToMallat (Real *mallat)
 int i_1 j_1 k_2int *lowHsize = new int [nsteps];
 int *lowVsize = new int [nsteps];
 lowKsize[nsteps-1] = (hsize+1)/2:
 lowVsize[nsteps-1] = [vslze+1]/2:
 for (i = nsteps-2; i > = 0; i--) (
   lowKsize[i] = (lowKsize[i+1]+1)/2lowVsize[i] = (lowVsize[i+1]+1)/2:
 // put linearized image in Mallat format
 // special case for LL subband
 for (j = 0; j < subbandVsize[0]; j_{++}]
   for (i = 0; i < subbandHsize[0]; i \rightarrowmallat[j*hsize+i] = subbandPtr[0][j*subbandHsize[0]+i];
 for (k = 0; k < nsteps; k \rightarrow) (
   for (j = 0; j < subbandVsize[k*3+1]; j++]
     for (i = 0; i < subbandHsize[k*3+1]; i++)
       mallat[j*hsize+(lowHsize[k]+i)] =subbandPtr[k*3+1][j*subbandHsize[k*3+1]+i];
   for (j = 0; j < subbandVsize[k*3+2]; j++)for (i = 0; i < subbandHsize[k*3+2]; i**]
       mallat(lowVsize[k]+1)*hsize+ii =subbandPtr[k*3+2][j*subbandHsize[k*3+2]+i];
```
for  $(i = 0; i <$  subbandVsize[0];  $(i+1)$ 

// move transformed image (in Mallat order) into linear array structure // special case for LL subband

int steps, int isSymmetric)

void WaveletTransform:: transform (Image \*image, Wavelet \*newWavelet,

wavelet->transform2d (image->value, temp, hsize, vsize, nsteps, symmetric);

wavelet->invert2d (temp, invertedImage->value, hsize, vsize,

nsteps, symmetric);

// clear out old info and set up subband pointers

void WaveletTransform::invert (Image \*invertedImage)

void WaveletTransform::mallatToLinear (Real \*mallat)

for  $(j = 0; j <$  subbandVsize[k\*3+3]; j++] for  $(i = 0; i <$  subbandHsize[ $k*3+3$ ];  $i++)$  $mailat[(lowVsize[k]+j)*hsize*(lowKsize[k]+i)] =$ 



## image.cpp

```
#include <stdio.h>
#include <stdlib.b>
#include <math.h>
*include "global.hh"
#include "image.hh"
// Create a blank image with width=hsize, height=vsize
// If hsize is unspecified, creates an image with width=0, height=0
// If vsize is unspecified, creates a square image with width =
\frac{1}{2}heicht = hsizeImage:: Image (int new hsize, int new vsize) : hsize(new_hsize),
  vsize(new vsize)
 if (hsize == -1)heiza = veiza = 0+if (vsize == -1)
  vsize = bisizevalue = new Real [hsize*vsize];
 if (value == NULL)
   error ("Can't allocate memory for image of size %d'by %d\n",
        hsize, vsize);
// Conv constructor
Image:: Image (const Image& image)
  int i;
  hsize = image.hsizevsize = image, vsize.value = new Real (hsize*vsize):
  if (value == NULL)
    error ('Can't allocate memory for image of size %d by %d\n',
         hsize, vsize);
  for (i = 0: i < hsize*vsize: i++)value[i] = image.value[i];// Loads a raw image of size hsize by vsize from the specified file
// The file is assumed to contain an image in raw byte format
Image:: Image (const char "filename, int new hsize, int new vsize) :
 hsize(new hsize), vsize(new vsize)
 if (hsize = -1)
  hsize = vsize = 0If (vsize == -11)
  vsize = hisizevalue = new Real (hsize*vsize);
 if (value == NULL)error ("Can't allocate memory for image of size %d by %d\n",
        hsize, vsize);
  loadRaw (filename);
```

```
// Loads a PGM image from a file. Sets hsize, vsize.
Image::Image (const char *filename)
 vsize = hsize = 0:
 value = NULLloadPGM (filename);
// Destructor
Image: "Image 1)hsize = vs1ze = -1;
  delete [] value;
// Assignment operator
Image & Image: : operator= (const Image& image)
 delete il value:
 hsize = image.bsizevsize = image.vsize;Value = new Real (hsize*vsizel:
 for \lim f i = 0: i < hsize*vsize: i++1
  value[i] = image.value[i];return *this.
// Loads an image from the specified file. The file is assumed to
// contain an image in raw byte format of size hsize by vsize
void Image:: loadRaw (const char *filename)
  FILE *infile;
  unsigned char *buffer:
  int i:
  infile = fopen (filename, "rb"):
  if (infile == NULL)
    error ("Unable to open file %s\n", filename);
  buffer = new unsigned char [hsize * vsize];
  if (fread (buffer, hsize*vsize, sizeof(unsigned char), infile) != 1)
      error ('Read < %d chars when loading file %s\n', hsize*vsize, filename);
  for (i = 0; i < hsize*vsize; i++)value[i] = (Real)buffer[i];delete il buffar:
  fclose (infile);
```
*II* **Saves an image to the specified file. The image is written in**  *II* raw byte format.

```
void Image: :saveRaw (const char *filename) 
(
```
FILE \*outfile: **unsigned char \*buffer; int i:** 

 $outfile = from (filename, 'wb+');$ if (outfile == NULL) error ('Unable to open file %s\n', filename);

buffer = new unsigned char [hsize\*vsize];

for (i = 0; i **< hsize\*vsize; i++)**  buffer[i] = realToChar(value[i]);

**fwrite (buffer, hsize\*vsize, 1, oucfile);** 

```
delete (1 buffer;
fclose [outfile);
```
*II* Private stuff to load pgms/ppms

void Image: : PGMSkipComments (FILE\* infile, unsigned char\* ch)

```
while ( l * ch = 1' * l ) (
   while \{\text{`ch} \equiv \langle \text{`h'} \rangle \} (\text{`ch} = \text{fgetc} (infile);
   while (*ch < ' ' ) ( *ch = fgetc (infile);
```

```
II Comment(s)
```
(

(

*II* **Get a number from a pgm file header, skipping comments etc.** 

unsigned int Image, ,PGMGetVal (FILE' infile)

```
unsigned int tmp; 
unsigned char chi 
do ( ch = fgetc(infile); ) while ((ch <= '') && (ch != '#'));
PGMSkipComments(infile, &ch);
ungetc(ch, infile); 
if (fsoanf(intile, "su", \&tmp) != 1) (
 printf("%s\n", "Error parsing file!");
 ext(1);
return (tmp) ;
```
*II* Loads a binary [PSI PGM image from the specified file. Sets hsize and *II* **vsize to the correct values for the file.** 

**void Image: :loadPGM (canst char \*filename)**  (

**FILE\* infile: unsigned char eh** *' ;* 

**infile = foper. (filename, - rb -);**   $if$   $\{infile == \text{NULL}\}$ 

*II* Look for type indicator while  $((ch \{p'p')\} \& (ch \{p'q'q')\})$  (  $ch = fgetc (infile)$ ; ) PGMSkipComments(infile, &ch); char ftype = fgetc(infile); *II* get type, S or 6 *II* **Look for x size, y size, max grey level**  int  $xsize = (int)PGMGetVal(intile);$  $int ysize = (int)PGMGetVal(intile);$ int maxg = (int)PGMGetVal(infile); *II* **Do some consistency checks**  if ( (hsize  $\leq 0$ ) && (vsize  $\leq 0$ ) ) ( **resize (xsize, ysize);**  if (value == NULL) error ('Can't allocate memory for image of *size* %d by %d\n', **hsize, vsize);**  else ( if ((xsize != hsize) || (ysize != vsize)) ( **error ( -File dimensions conflict with image settings\n- );**   $\mathbf{1}$  $\mathcal{N}$ if  $If type == '5'$ ) ( printf['Pile %s is of type PGM, is %d x %d with max gray level %d\n ', **filename , hsize, vsize, maxg);**  PGMLoadData(infile, filename); if (ftype ==  $'6'$ ) ( printf('Pile %s is of type PPM, *is* %d <sup>x</sup>%d with max gray level %d\n ', **filename, hsize, vsize, maxg);**  error('Attempt to load a PPM as a PGM\n');  $fclose(intfile)$ *II* Loads the data segment of a PGM image from the specified file . **void Image::PGMLoadOata (FILE \*infile, const char \*filenarne)**  ( **unsigned char \*buffer; int i; buffer new unsigned char Ihsize \* vsizel; long fp = -l\*hsize\*vsize:**  fseek(infile, fp, SEEK\_END); if (tread (buffer, hsize\*vsize, sizeot(unsigned char), infile) != 1) **error ( - Read < %d chars when loading file %s\n-, hsize\*vsize, filename);**  for  $(i = 0; j < k$  hsize\*vsize;  $(i++)$  $value[i] = (Real)buffer[i];$ delete II buffer;

error ('Unable to open file %s\n', filename);

*1\* ------ II* Saves an image to the specified file. The image is written in